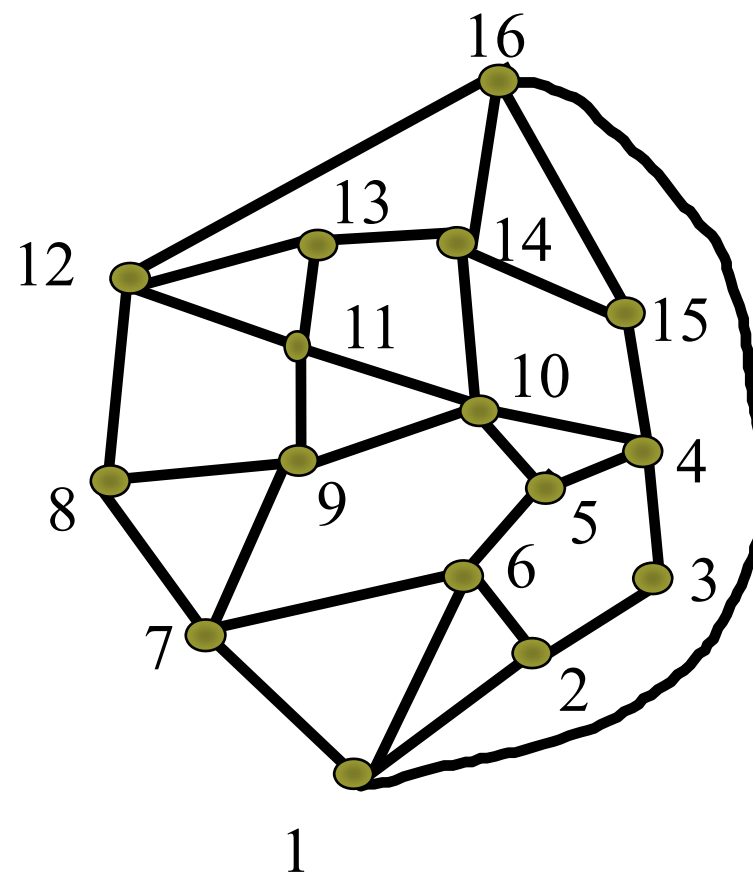


# Various Orders and Drawings of Plane Graphs

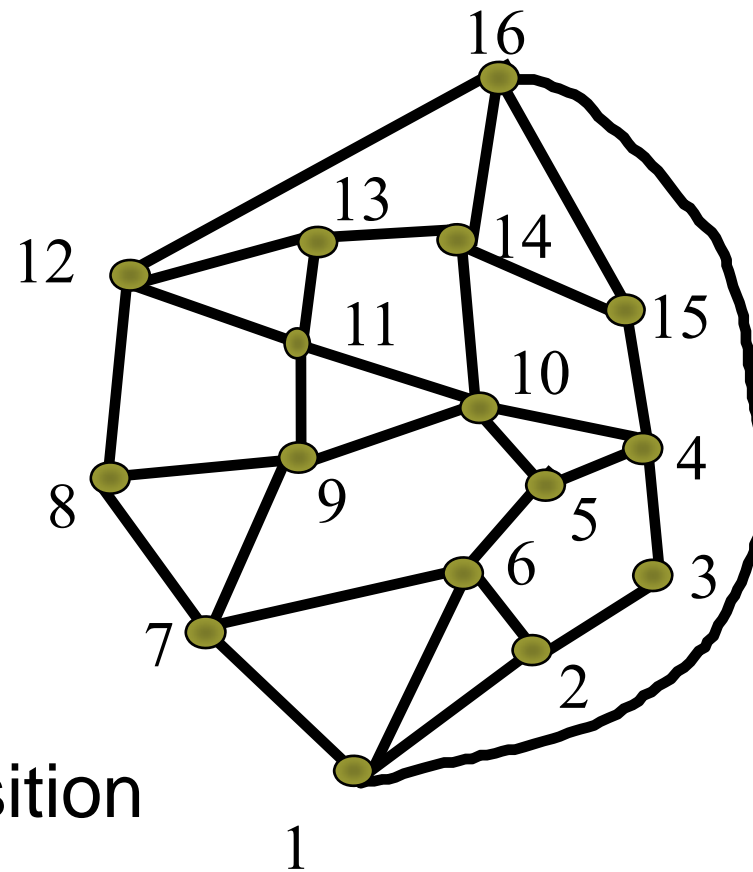
Takao Nishizeki  
Tohoku University

# Vertex ordering



# Vertex ordering

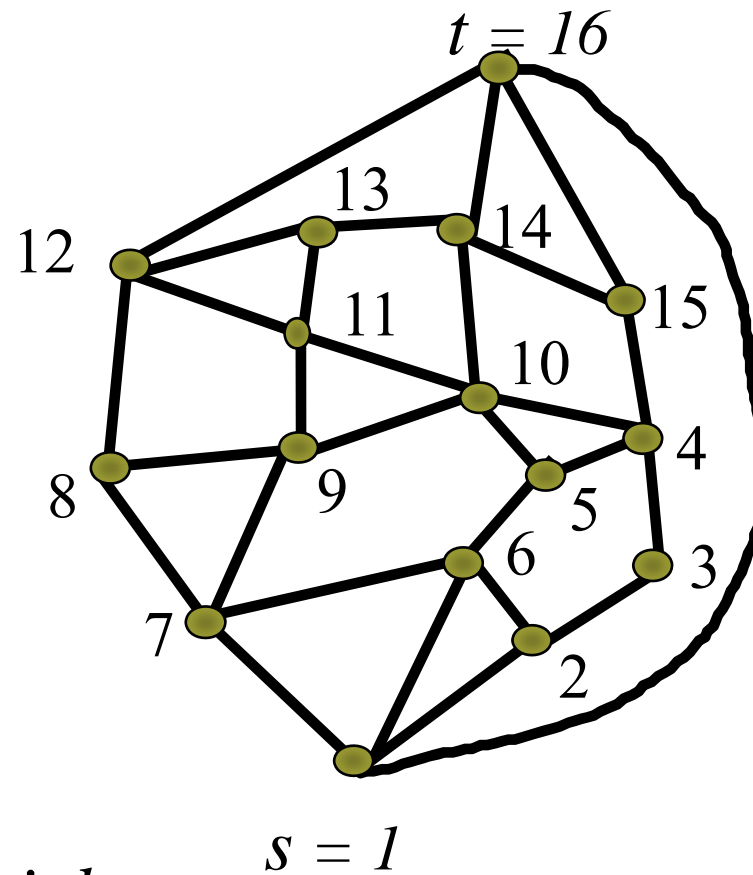
- *st*-numbering
- Canonical ordering
- 4-canonical ordering
- Canonical decomposition
- 4-canonical decomposition



## Vertex ordering

- *st*-numbering
- Canonical ordering
- 4-canonical ordering
- Canonical decomposition
- 4-canonical decomposition

# *st*-numbering

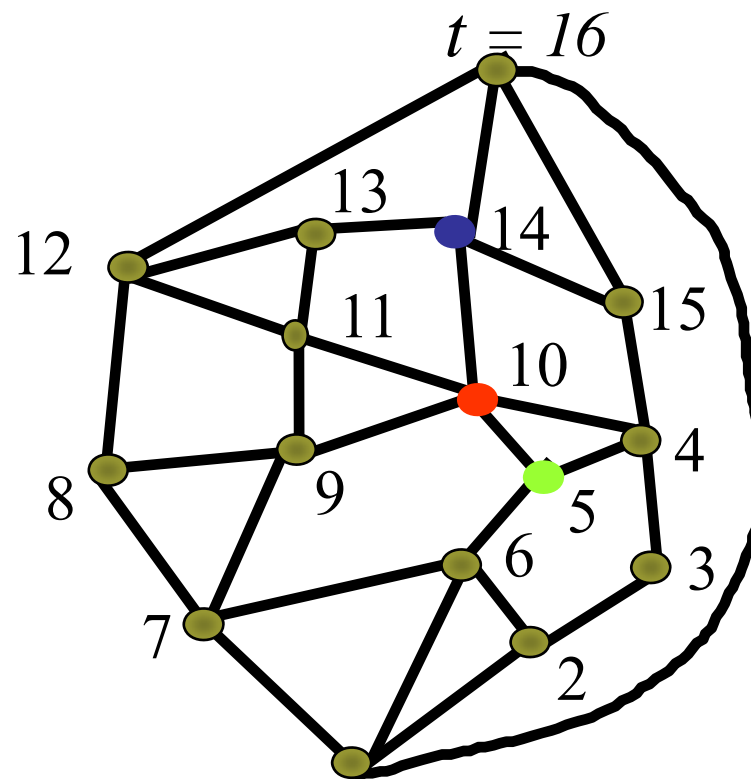


$i \neq s, t$  has two neighbors  $j, k$

$$j < i < k.$$

# *st*-numbering

$$i = 10$$



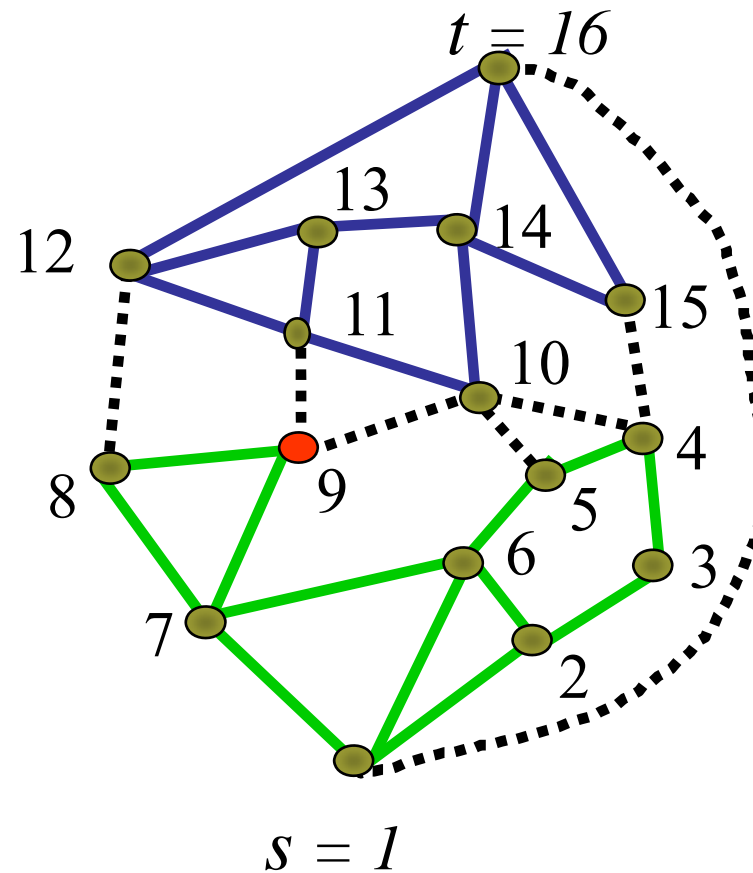
$$s = 1$$

$i \neq s, t$  has two neighbors  $j, k$

$$j < i < k.$$

# *st*-numbering

$$i = 9$$



For any  $i$ , both vertices  $\{1, 2, \dots, i\}$  and  $\{i + 1, i + 2, \dots, n\}$  induce connected subgraphs.

# Application of *st*-numbering

Planarity testing

Visibility drawing

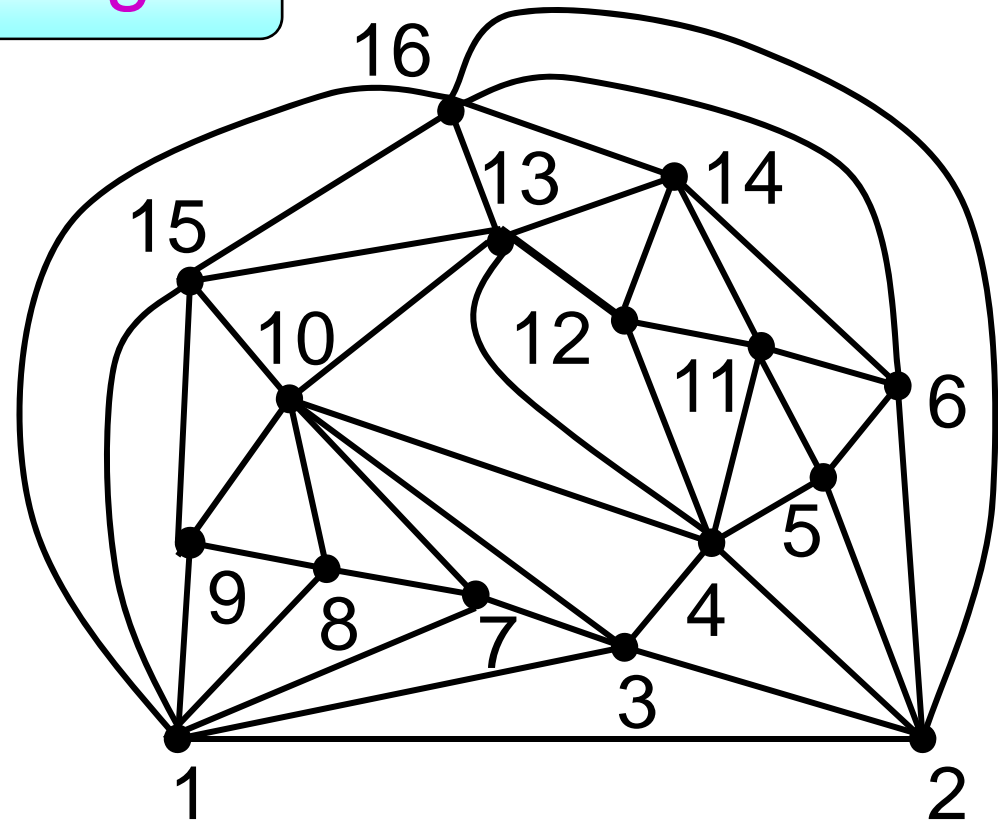
Internet routing



## Vertex ordering

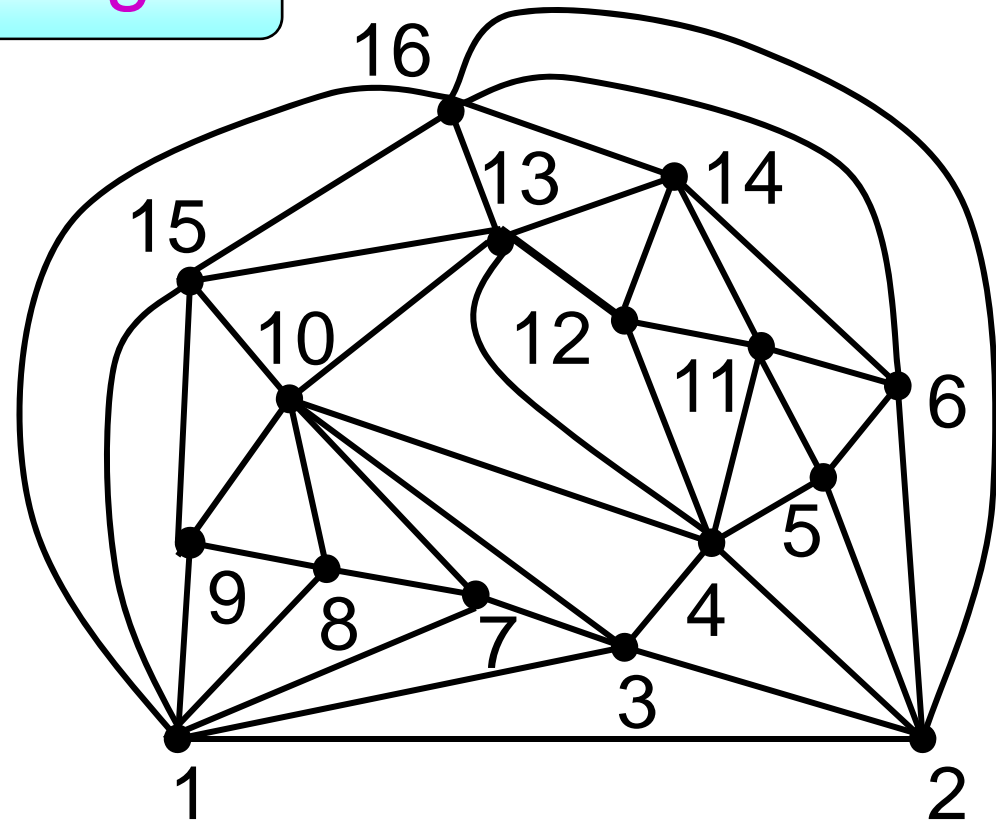
- *st*-numbering
- Canonical ordering
- 4-canonical ordering
- Canonical decomposition
- 4-canonical decomposition

# Canonical Ordering



Triangulated plane graph

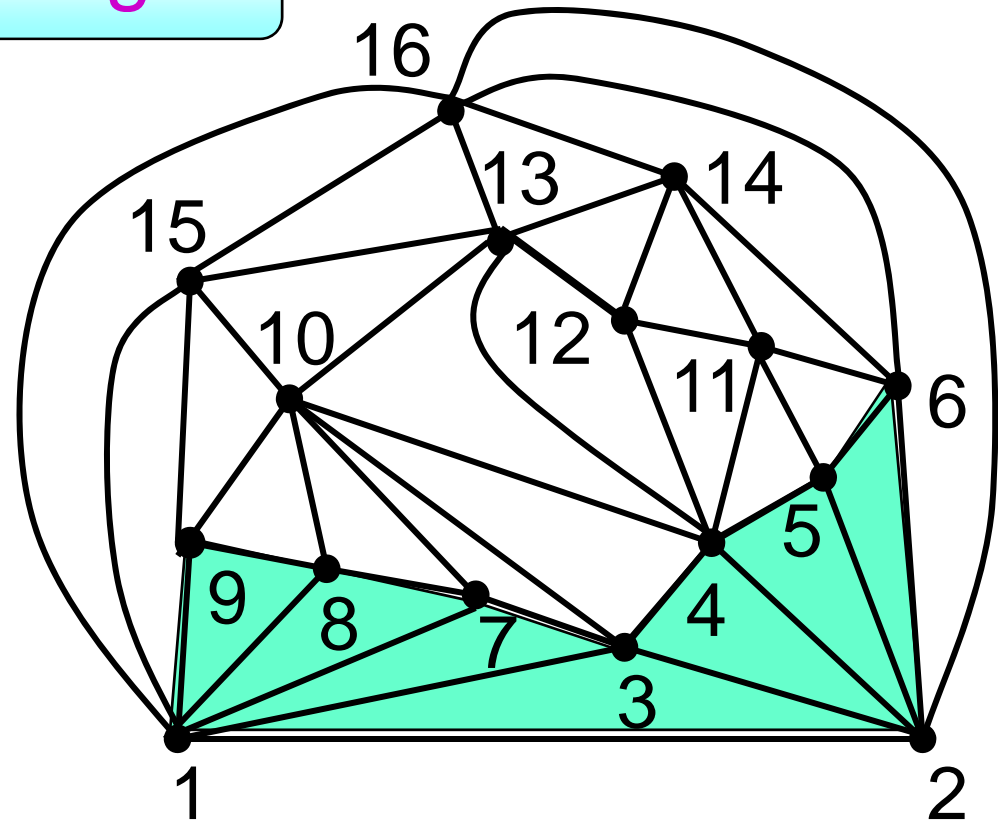
# Canonical Ordering



$G_k$ : subgraph of  $G$  induced by vertices  $1, 2, \dots, k$

# Canonical Ordering

$G_9$

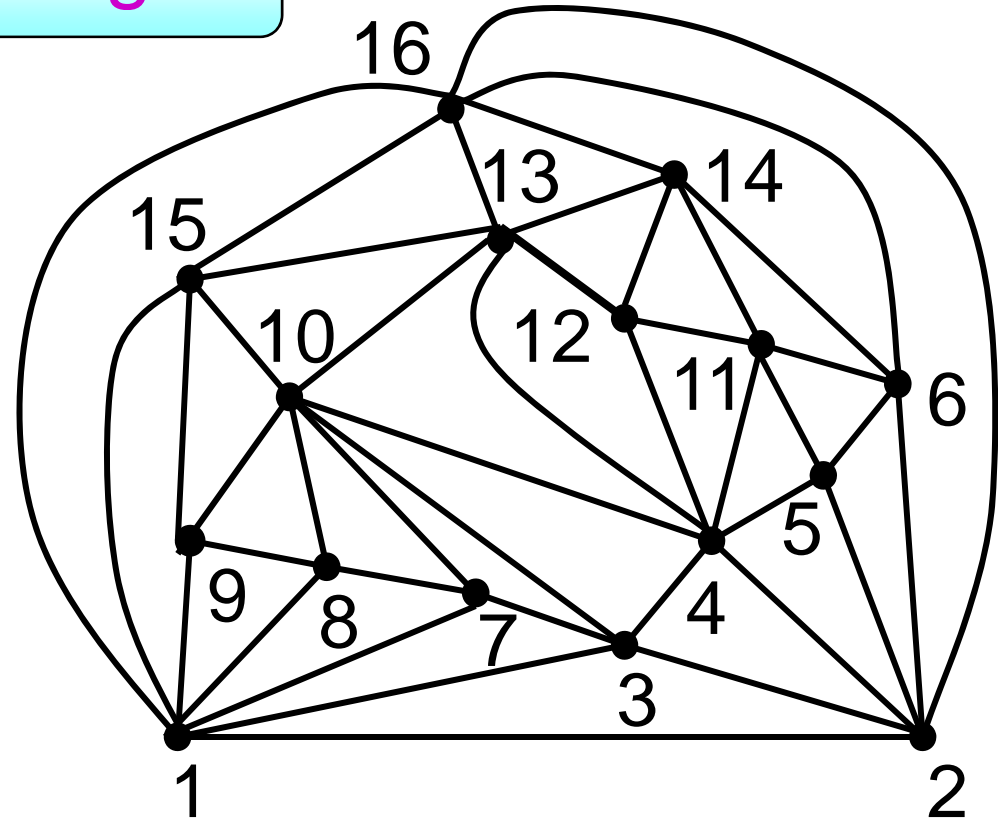


$G_k$ : subgraph of  $G$  induced by vertices  $1, 2, \dots, k$

# Canonical Ordering

For any  $k, 3 \leq k \leq n$

(co1)  $G_k$  is biconnected and internally triangulated



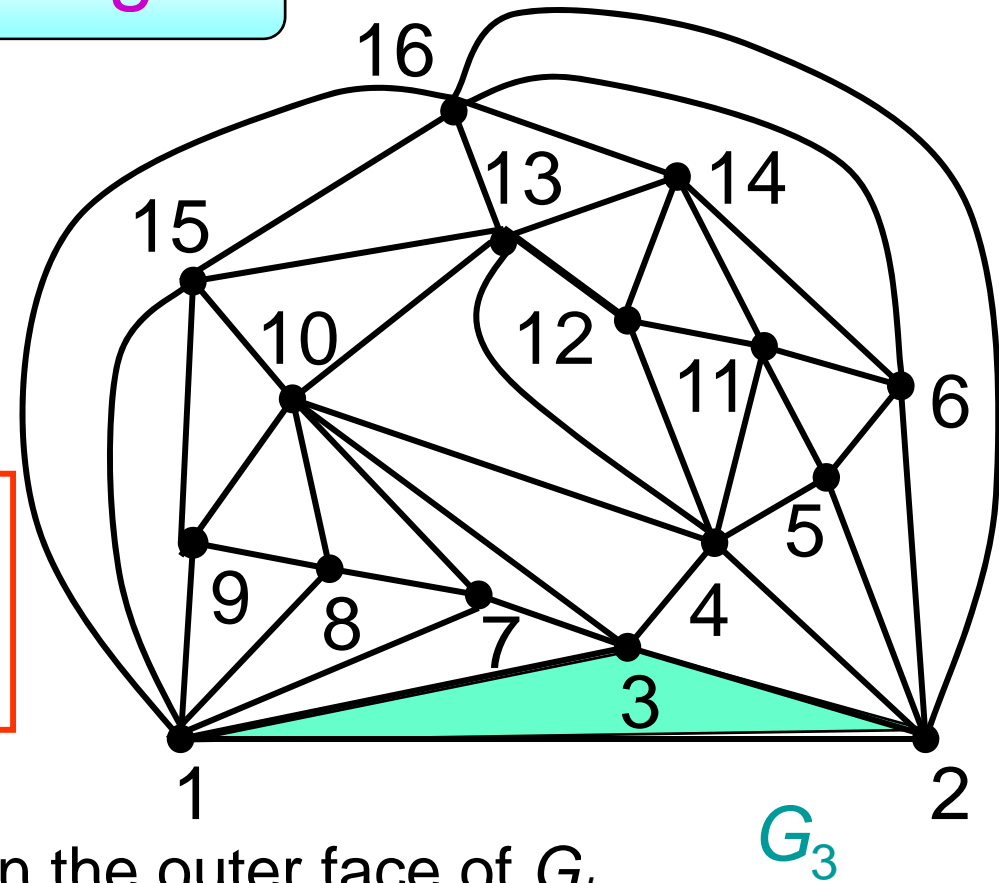
(co2) vertices 1 and 2 are on the outer face of  $G_k$

(co3) vertex  $k+1$  is on the outer face of  $G_k$  and the neighbor of  $k+1$  is consecutive on the outer cycle  $C_o(G_k)$ .

# Canonical Ordering

For any  $k, 3 \leq k \leq n$

(co1)  $G_k$  is biconnected and internally triangulated



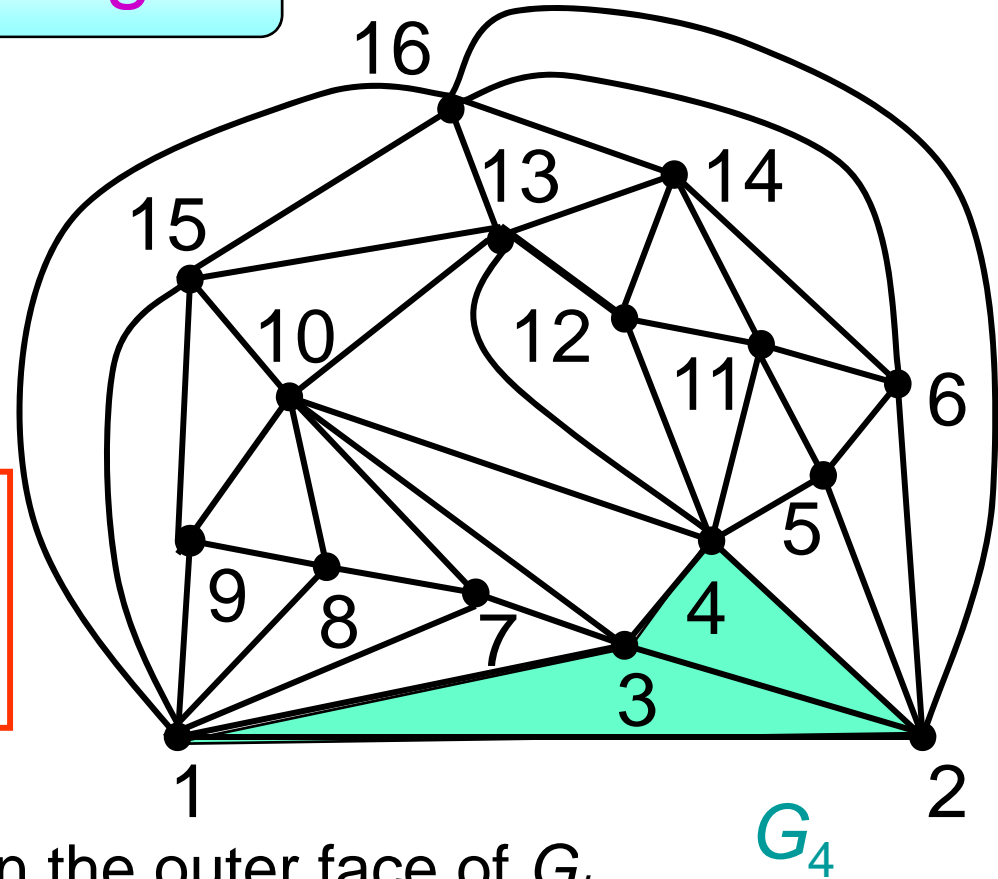
(co2) vertices 1 and 2 are on the outer face of  $G_k$

(co3) vertex  $k+1$  is on the outer face of  $G_k$  and the neighbor of  $k+1$  is consecutive on the outer cycle  $C_o(G_k)$ .

# Canonical Ordering

For any  $k, 3 \leq k \leq n$

(co1)  $G_k$  is biconnected and internally triangulated



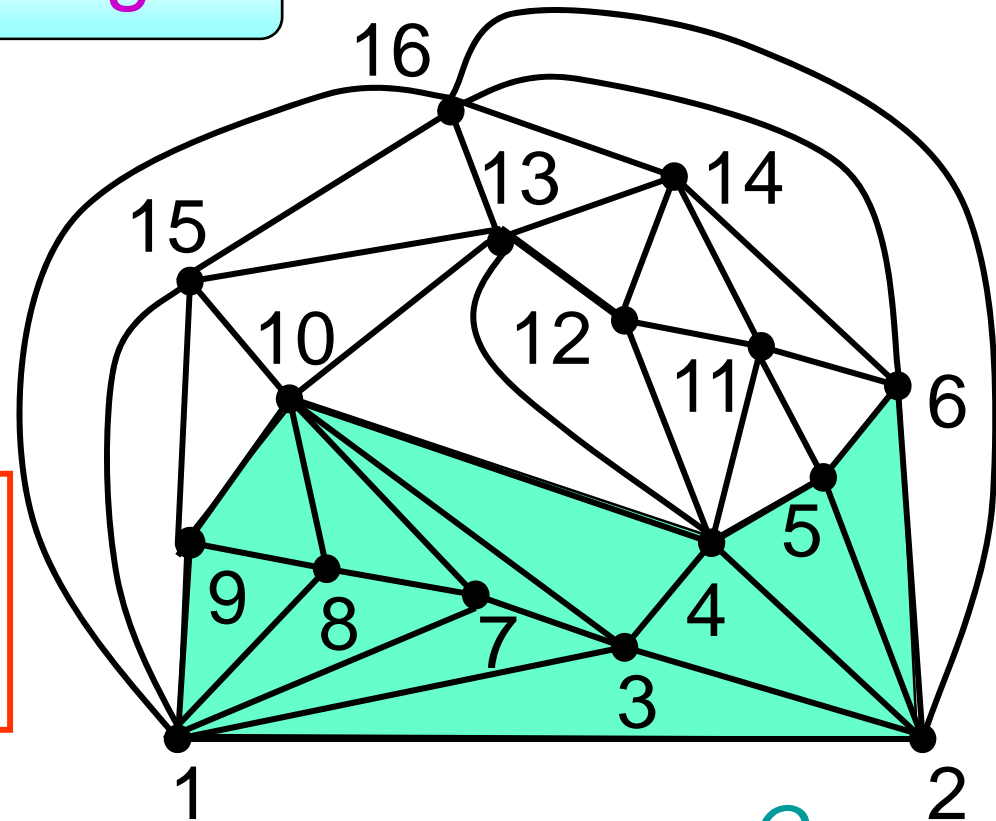
(co2) vertices 1 and 2 are on the outer face of  $G_k$

(co3) vertex  $k+1$  is on the outer face of  $G_k$  and the neighbor of  $k+1$  is consecutive on the outer cycle  $C_o(G_k)$ .

# Canonical Ordering

For any  $k, 3 \leq k \leq n$

(co1)  $G_k$  is biconnected and internally triangulated



(co2) vertices 1 and 2 are on the outer face of  $G_k$

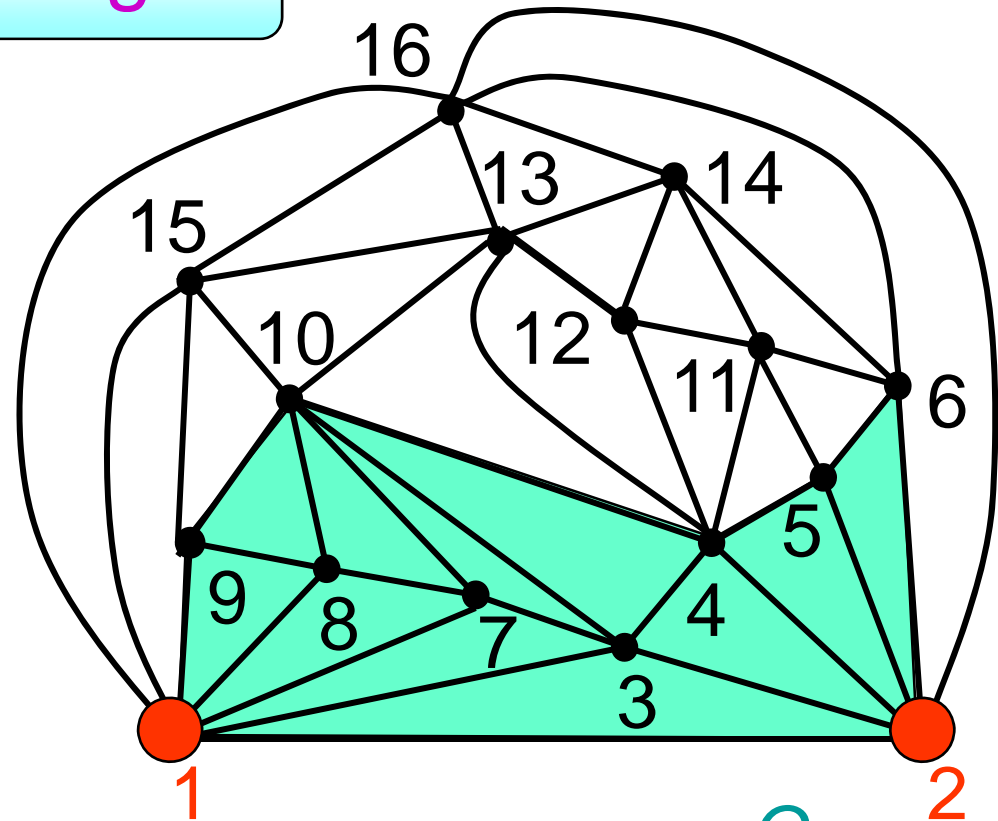
(co3) vertex  $k+1$  is on the outer face of  $G_k$  and the neighbor of  $k+1$  is consecutive on the outer cycle  $C_o(G_k)$ .



# Canonical Ordering

For any  $k, 3 \leq k \leq n$

(co1)  $G_k$  is biconnected and internally triangulated



(co2) vertices 1 and 2 are on the outer face of  $G_k$

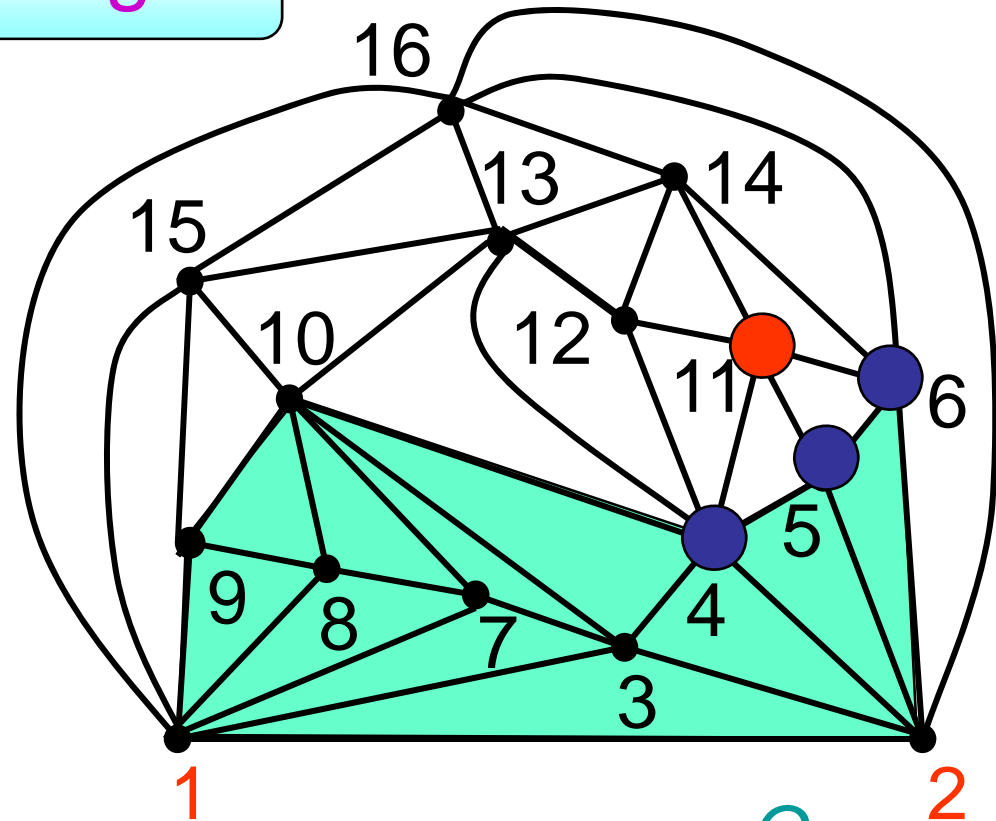
$G_{10}$

(co3) vertex  $k+1$  is on the outer face of  $G_k$  and the neighbor of  $k+1$  is consecutive on the outer cycle  $C_o(G_k)$ .

# Canonical Ordering

For any  $k, 3 \leq k \leq n$

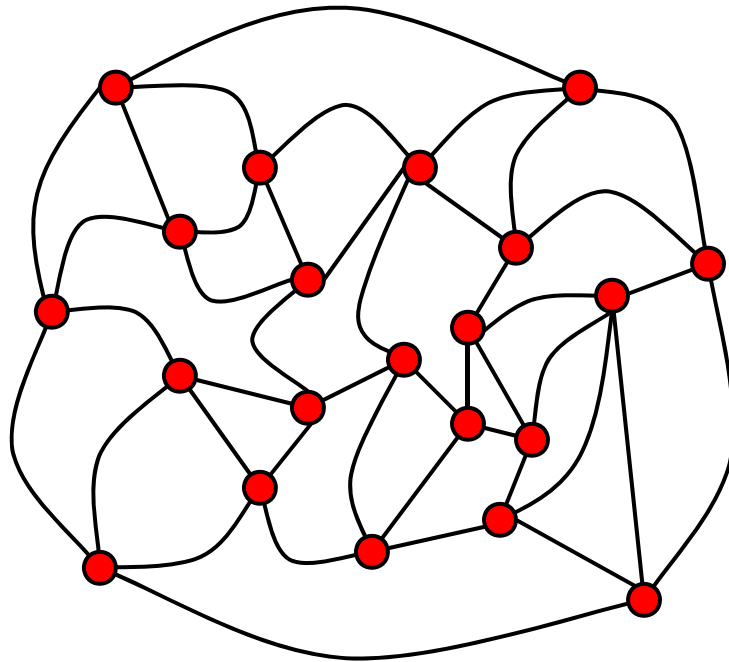
(co1)  $G_k$  is biconnected and internally triangulated



(co2) vertices 1 and 2 are on the outer face of  $G_k$

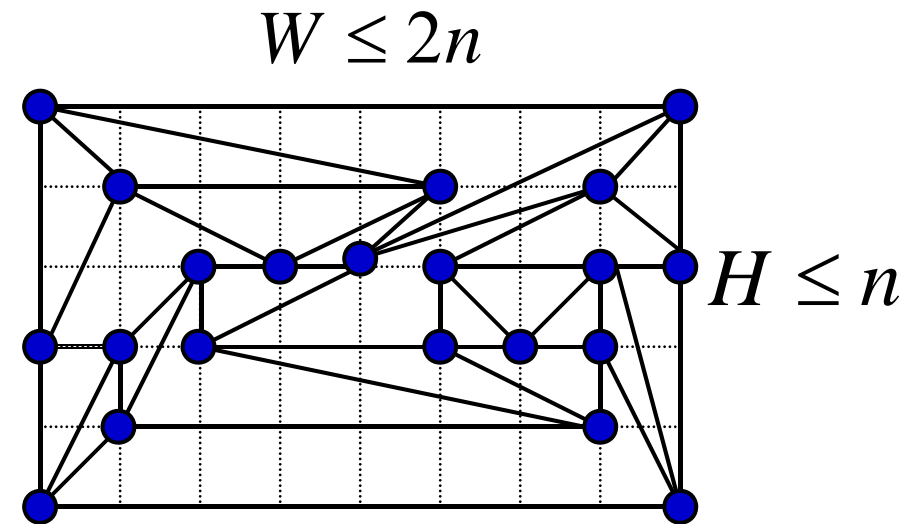
(co3) **vertex  $k+1$**  is on the outer face of  $G_k$  and the **neighbor of  $k+1$**  is consecutive on the outer cycle  $C_o(G_k)$ .

# Straight Line **Grid** Drawing



Plane graph

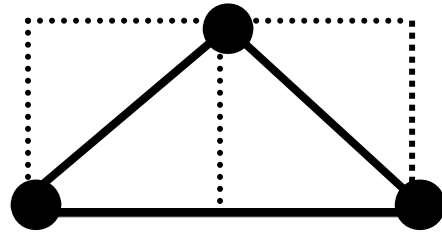
de Fraysseix *et al.* '90



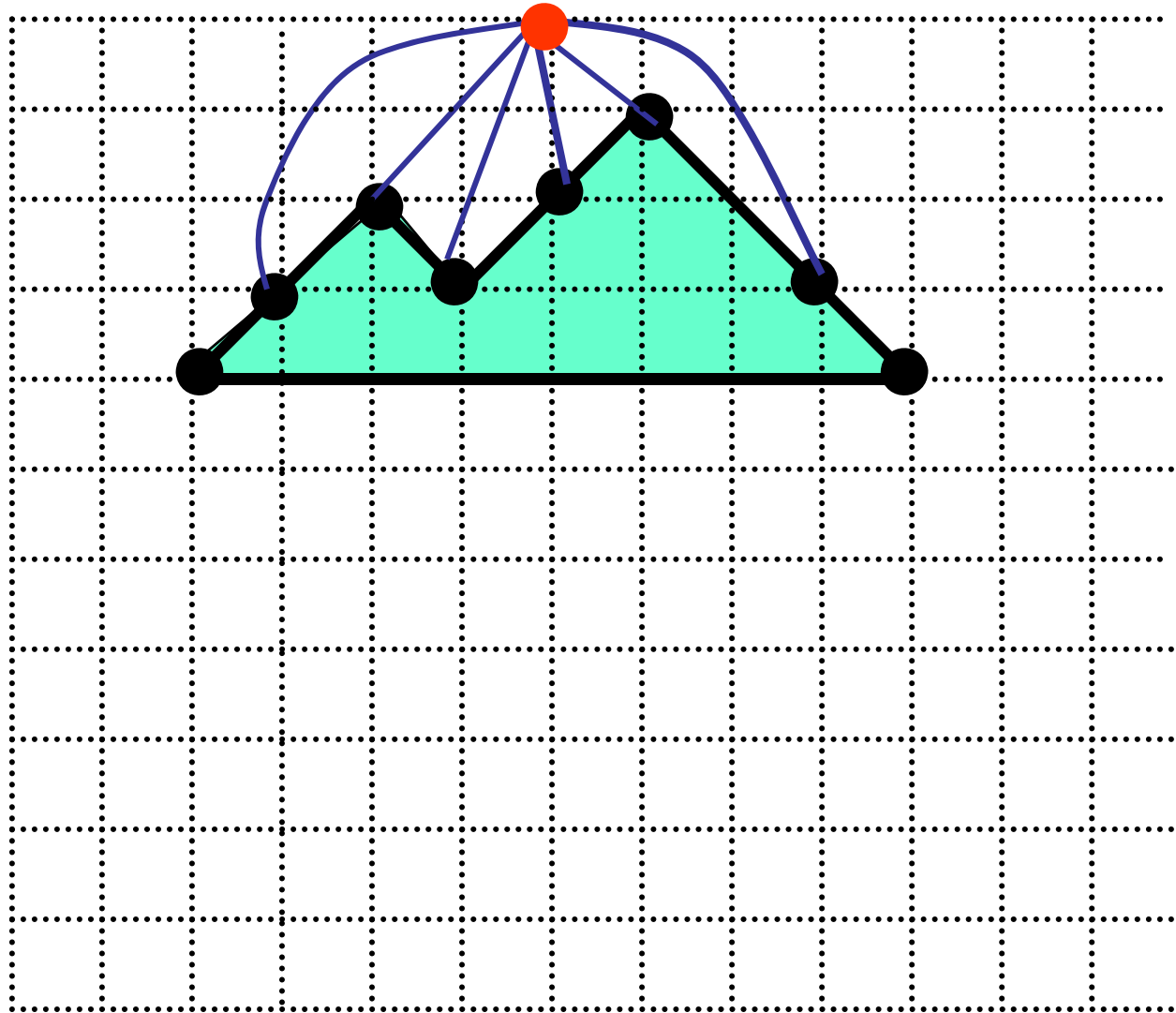
Straight line grid drawing.

$$W \times H \leq 2n^2$$

Initial drawing of  $G_3$



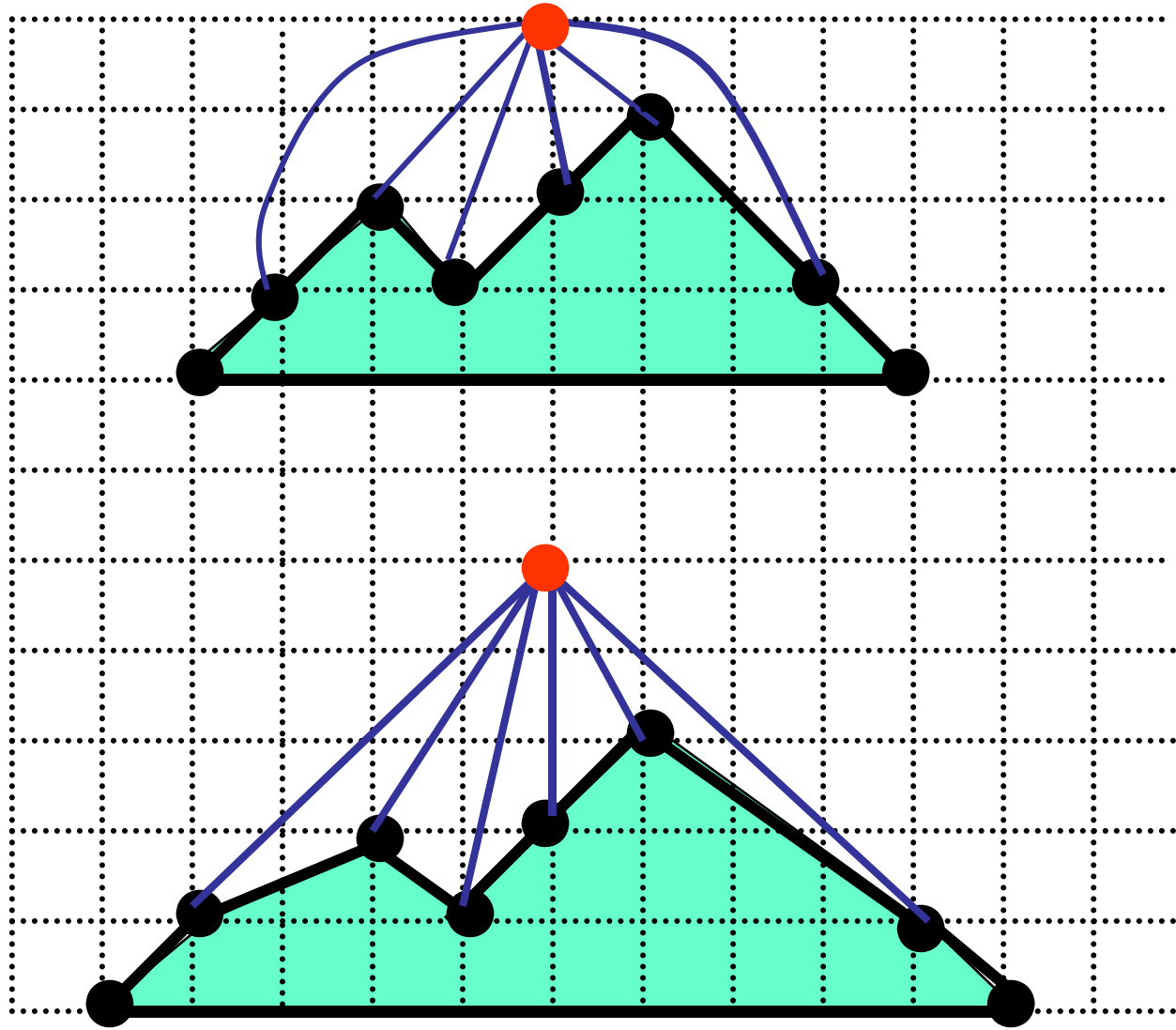
Install  $k + 1$



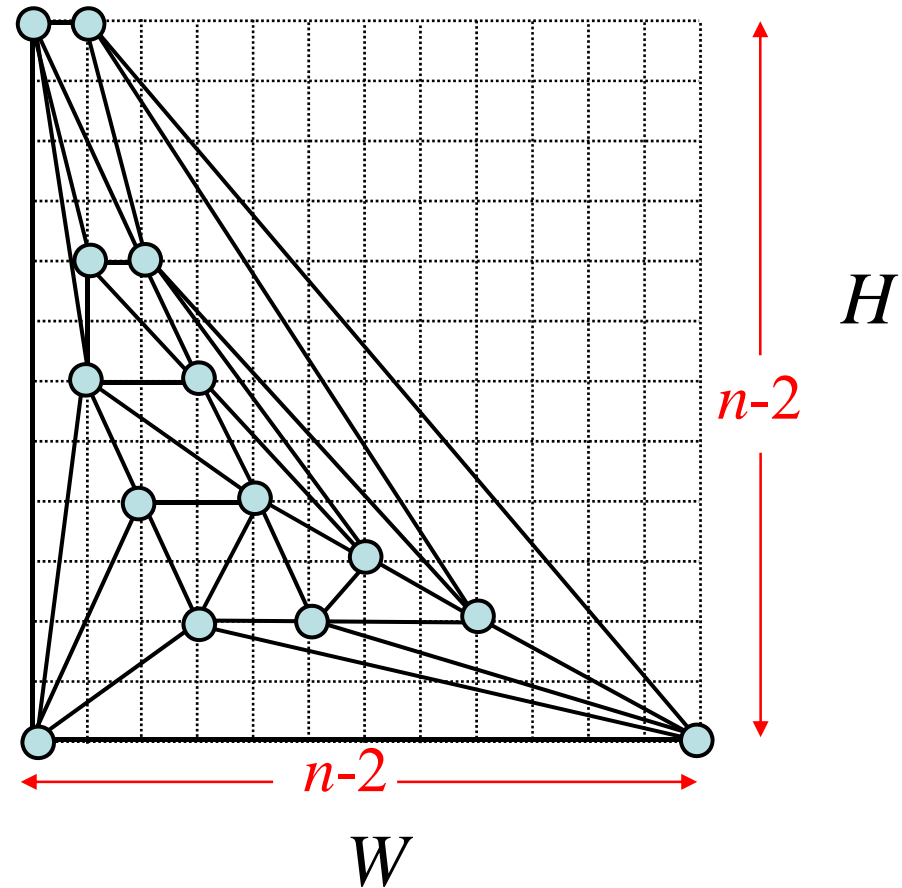
$G_k$

Shift method

Shift and install  $k + 1$



Schnyder '90



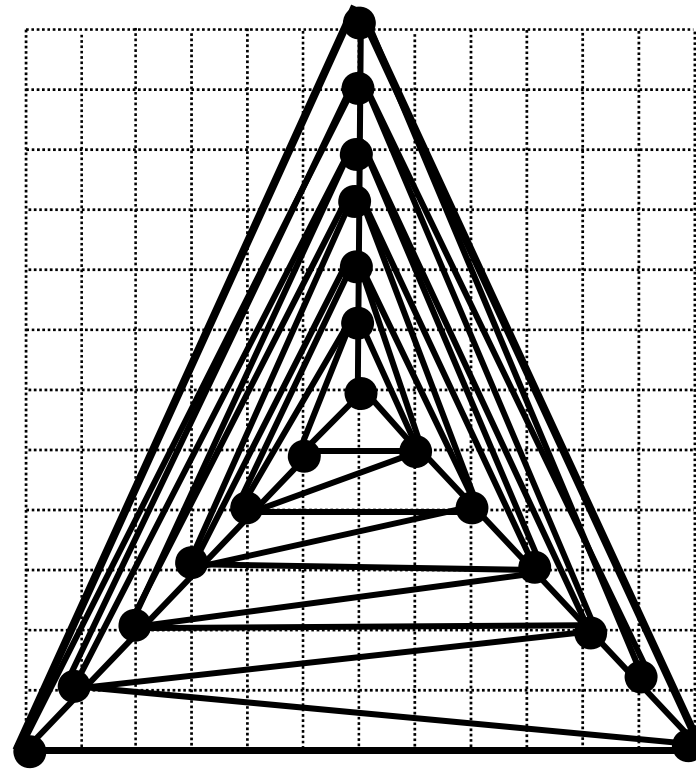
$$W \times H \leq n^2$$

Upper bound

What is the **minimum** size of a grid required for a straight line drawing?



# Lower Bound



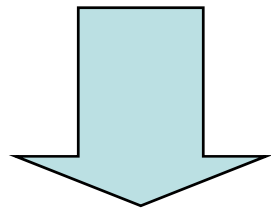
$$H \geq \frac{2}{3}n$$

$$W \geq \frac{2}{3}n$$

$$W \times H \geq \left(\frac{2}{3}n\right)^2$$

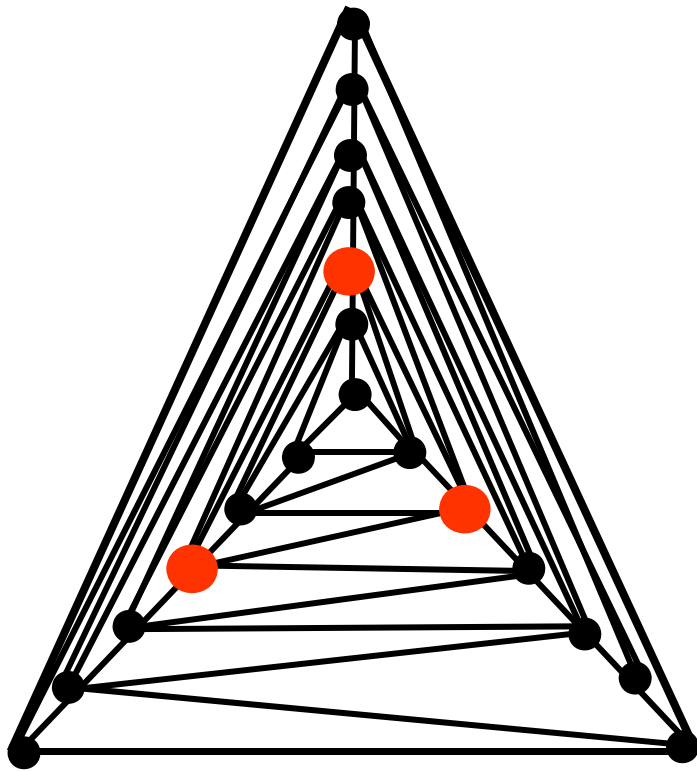
A **restricted class** of plane graphs may have more **compact** grid drawing.

Triangulated plane graph

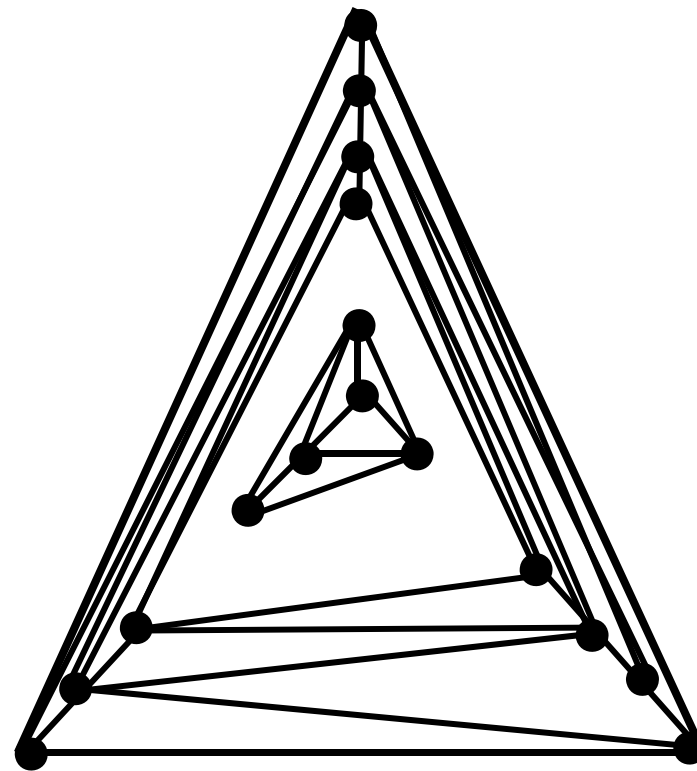


**3-connected graph**

4-connected ?



not 4-connected



disconnected

How much area is required for 4-connected  
plane graphs?

# Straight line grid drawing

Miura *et al.* '01

Input: 4-connected plane graph  $G$

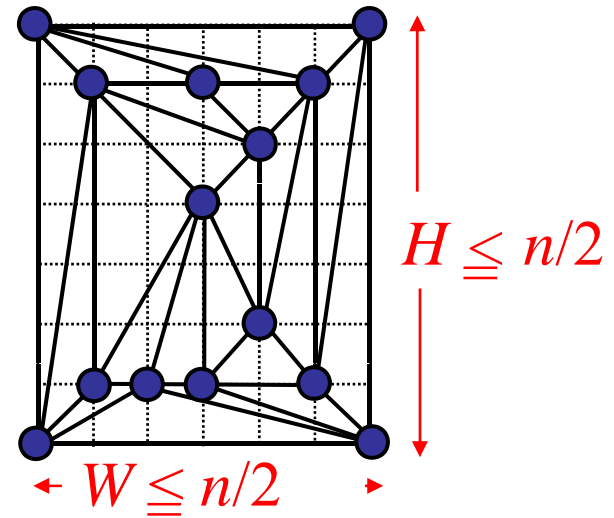
Output: a straight line grid drawing

Grid Size :

$$W, H \leq \frac{n}{2}$$

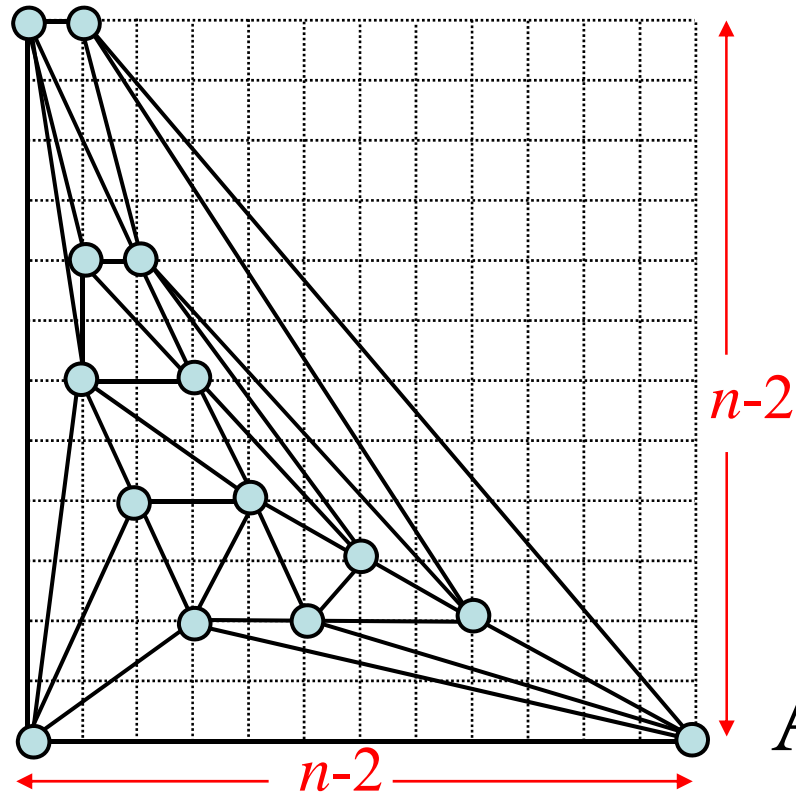
Area:

$$W \times H \leq \frac{n^2}{4}$$



Schnyder '90

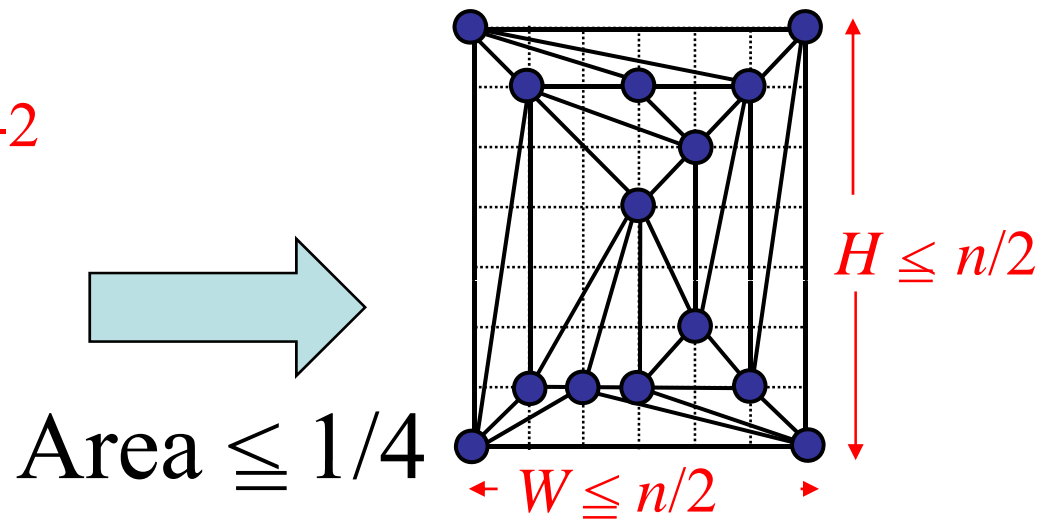
plane graph  $G$



$$\text{Area} \doteq n^2$$

Miura *et al.* '01

4-connected plane graph  $G$



$$\text{Area} \leq 1/4$$

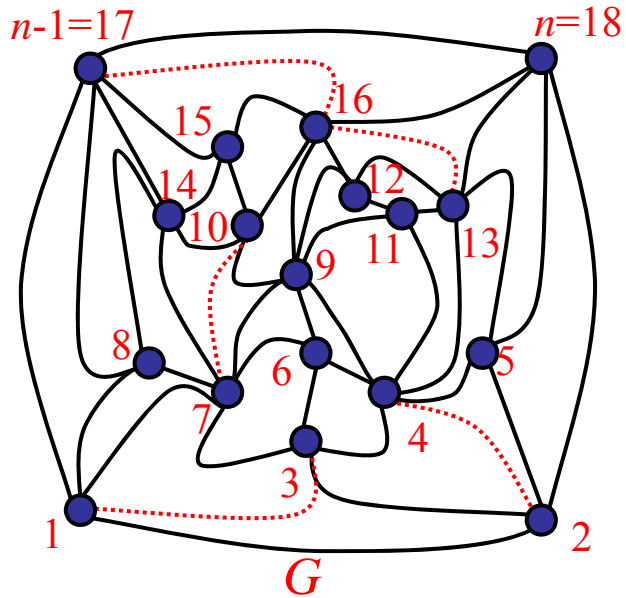
$$\text{Area} \leq n^2/4$$

## Vertex ordering

- *st*-numbering
- Canonical ordering
- 4-canonical ordering
- Canonical decomposition
- 4-canonical decomposition

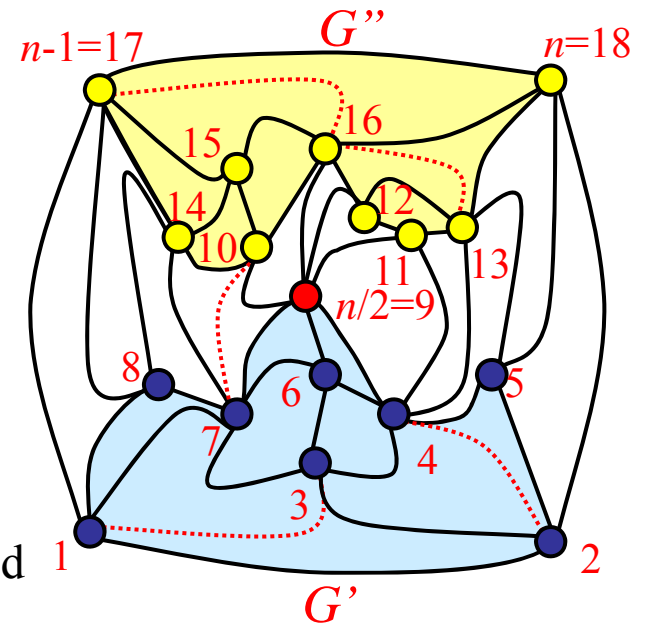
Triangulate all inner faces

Step1: find a 4-canonical ordering

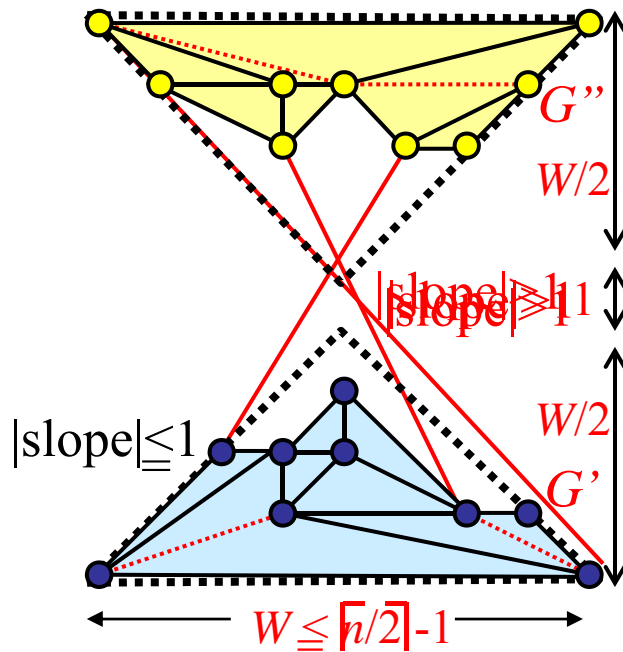
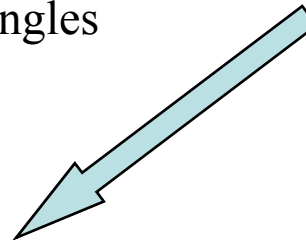


# Main idea

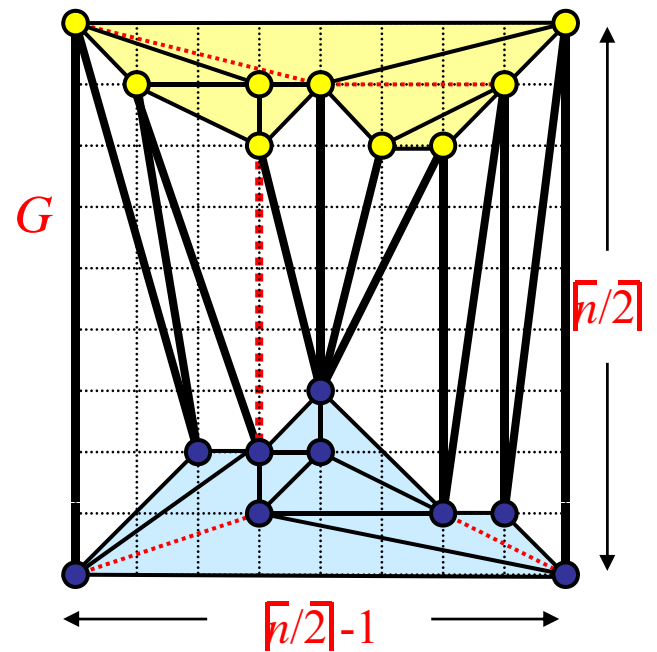
Step2: Divide  $G$  into two halves  $G'$  and  $G''$



Step3 and 4 : Draw  $G'$  and  $G''$  in isosceles right-angled triangles



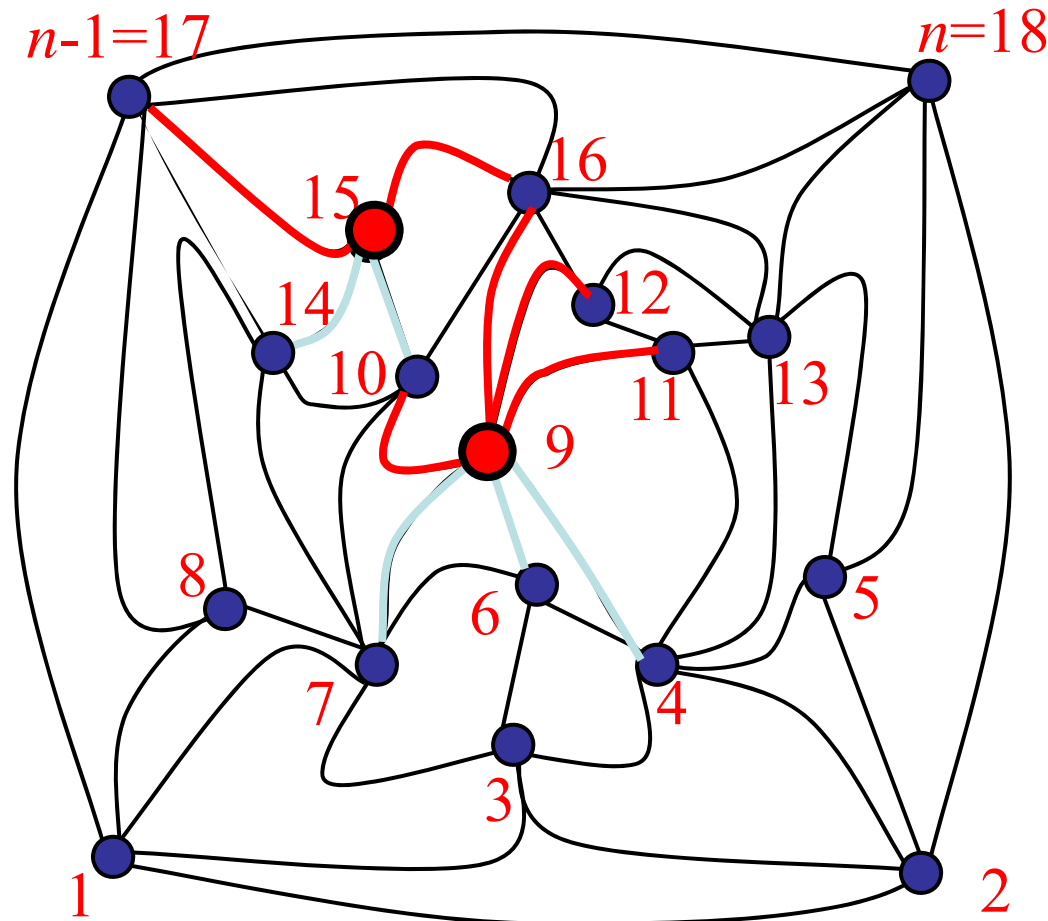
Step5: Combine the drawings of  $G'$  and  $G''$



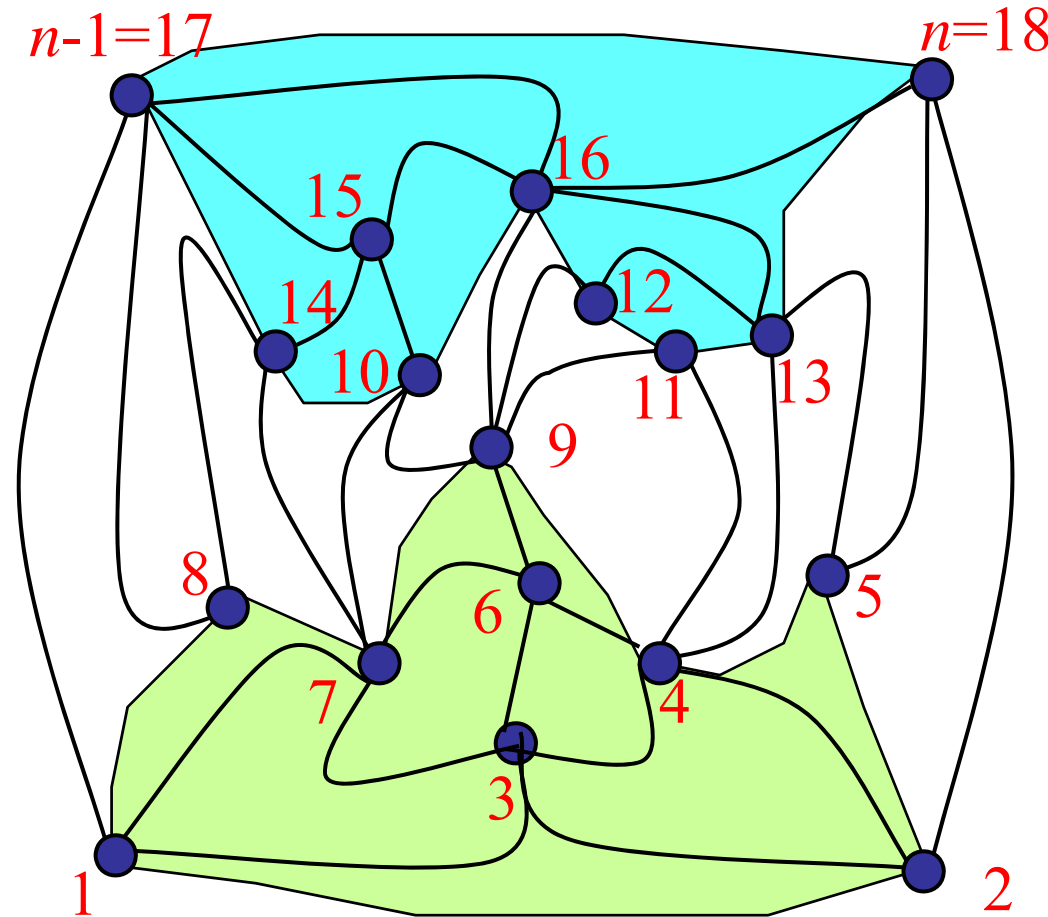


## 4-canonical ordering [KH97] (4-connected graph)

- (1) Edges  $(1,2)$  and  $(n,n-1)$  are on the outer face
- (2) For each vertex  $k$ ,  $3 < k < n-2$ , at least two neighbors have lower number and at least two neighbors have higher number.



# 4-canonical ordering [KH97] (4-connected graph)

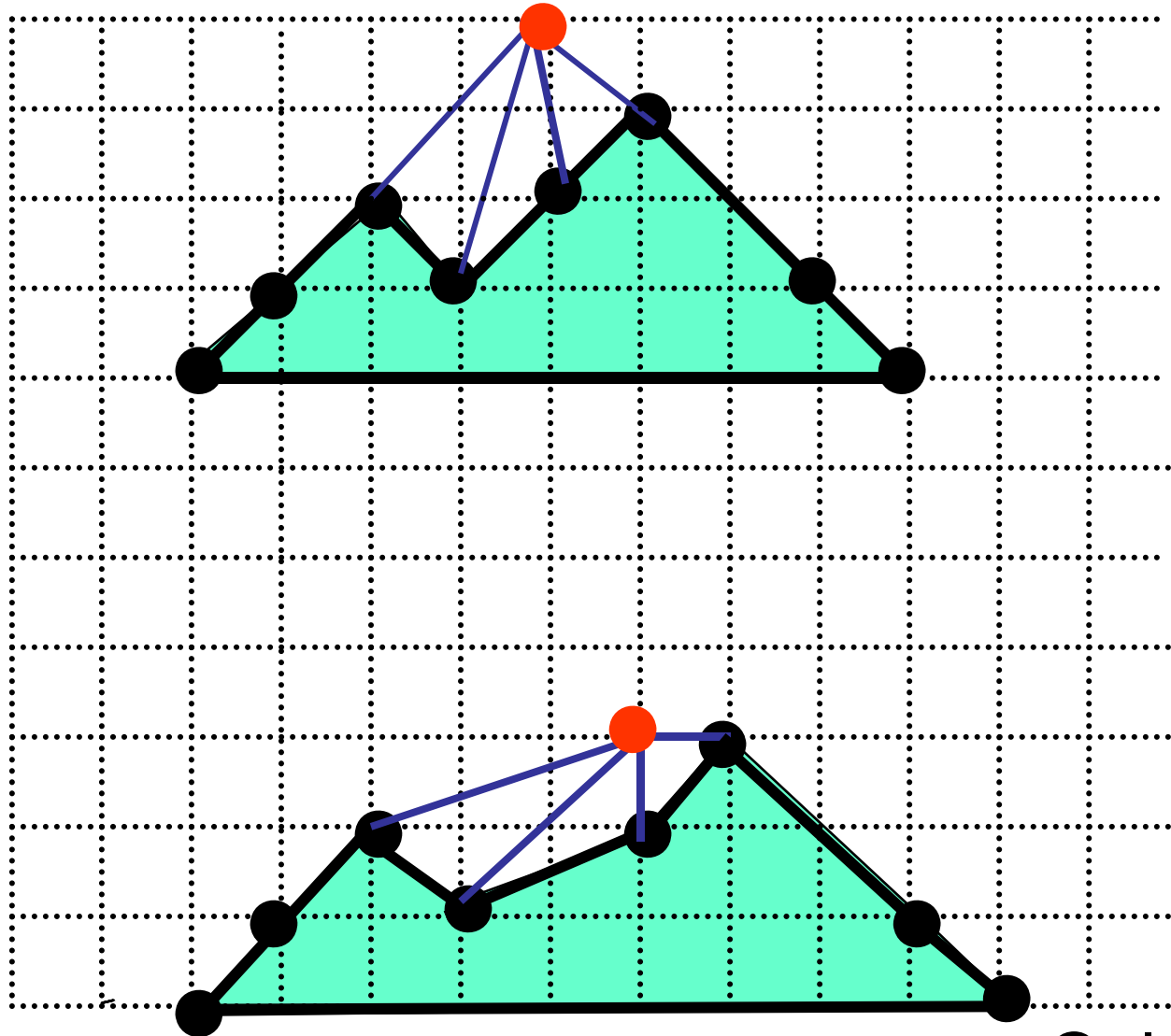


Generalization of  
an  $st$ -numbering

Both vertices  $\{1, 2, \dots, i\}$  and  $\{i + 1, i + 2, \dots, n\}$   
induce 2-connected subgraphs.

Shift method

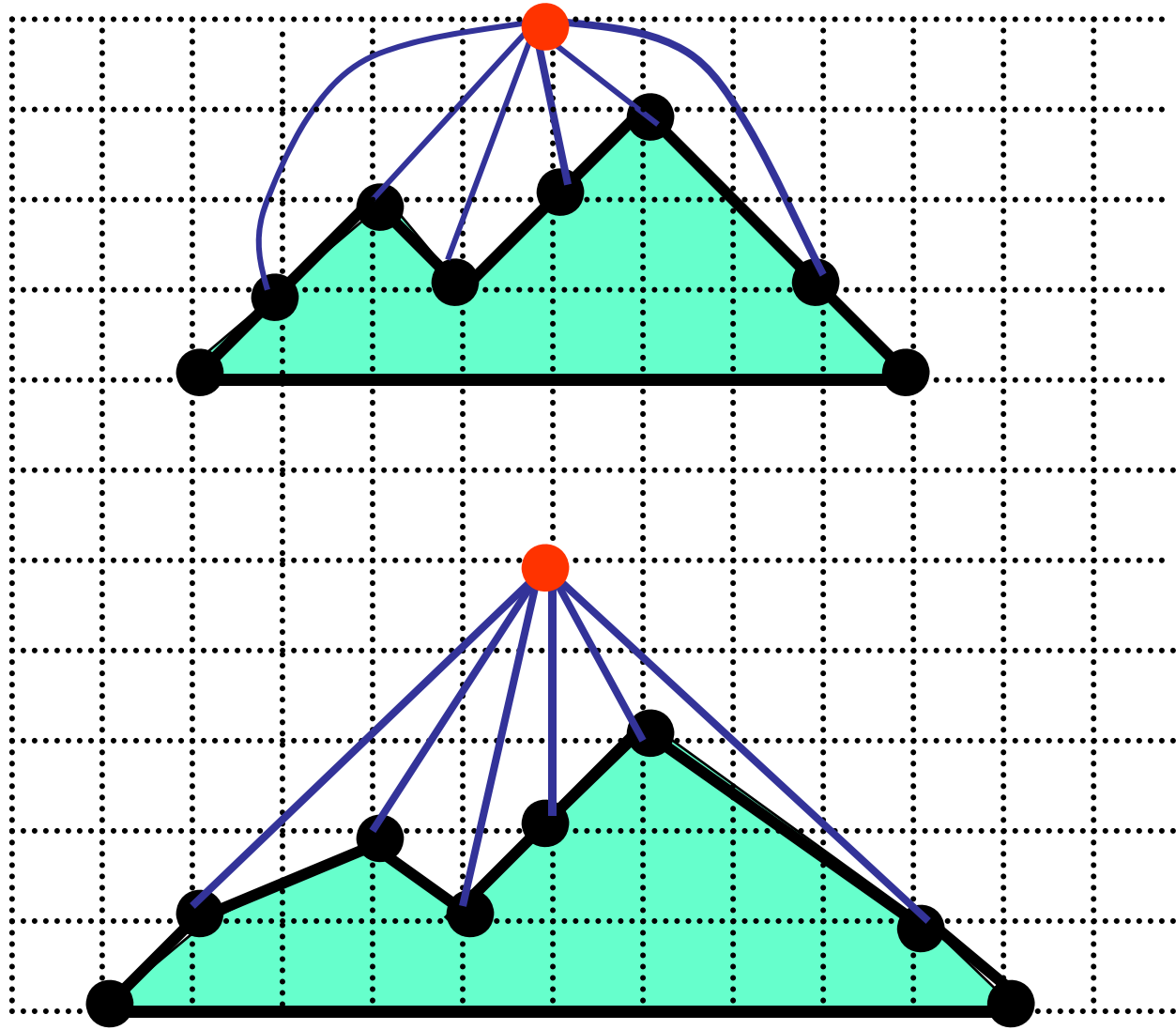
Shift and install  $k + 1$



Only one shift

Shift method

Shift and install  $k + 1$



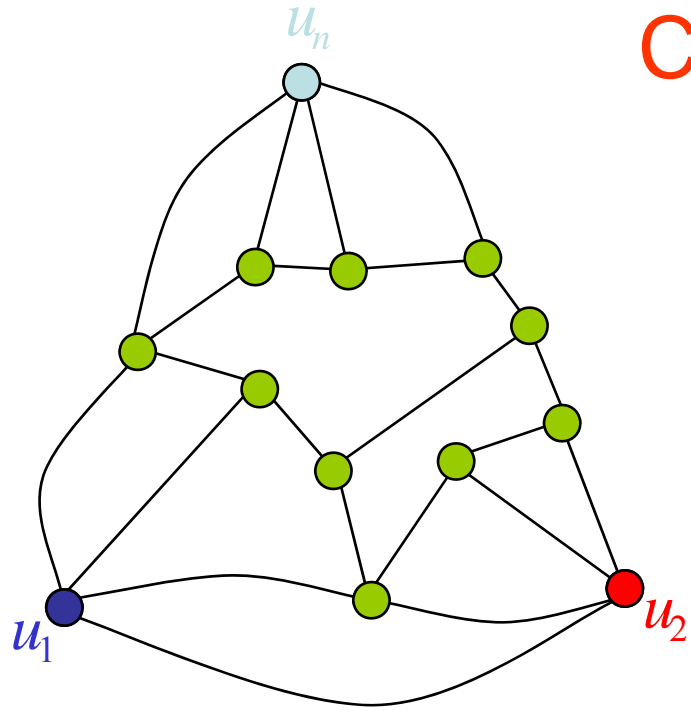
Graph is not triangulated

Is there any ordering?

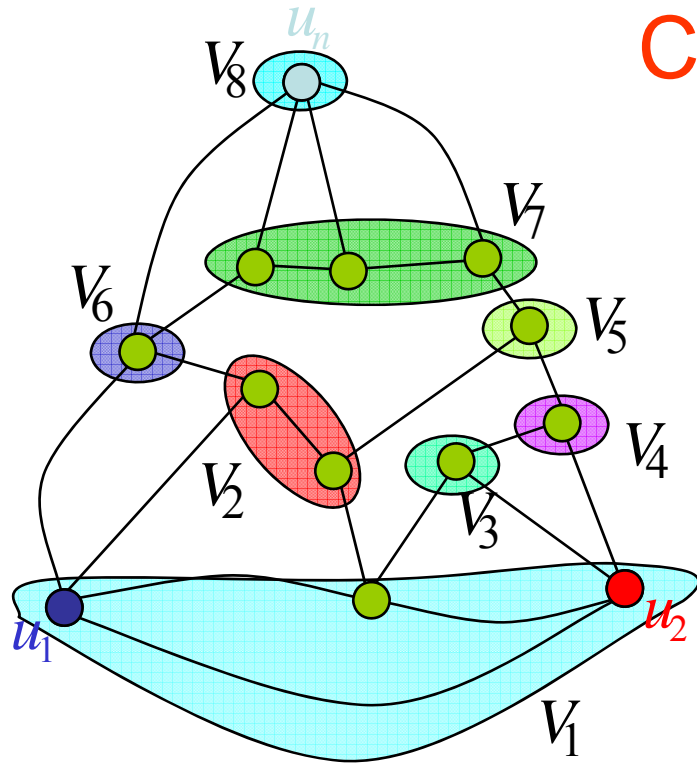
## Vertex ordering

- *st*-numbering
- Canonical ordering
- 4-canonical ordering
- Canonical decomposition
- 4-canonical decomposition

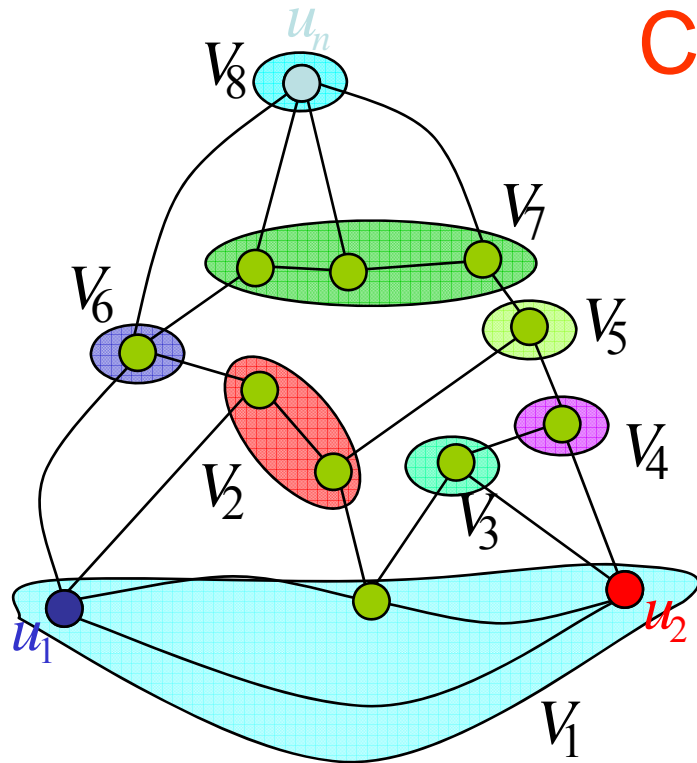
# Canonical Decomposition



# Canonical Decomposition





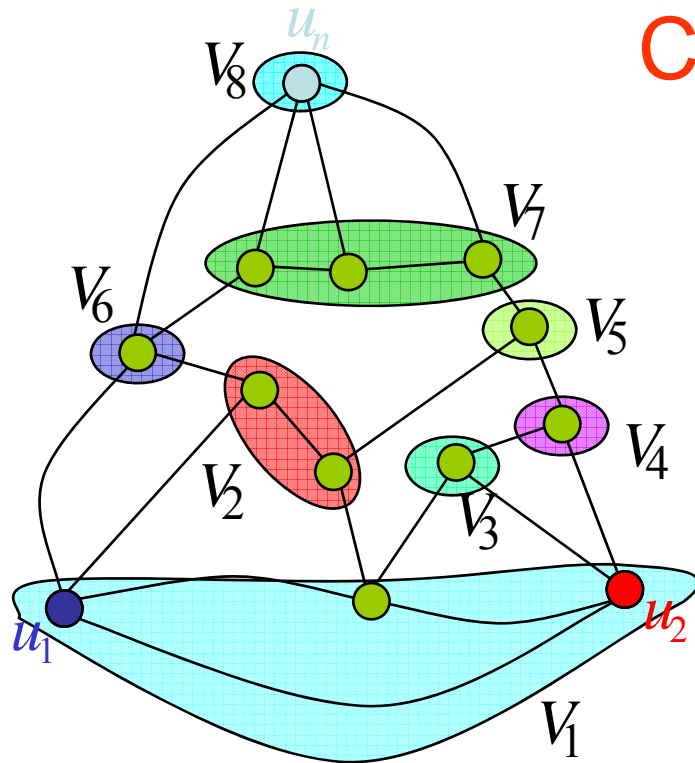


## Canonical Decomposition

(cd1)  $V_1$  is the set of all vertices on the inner face containing edge  $(u_1, u_2)$ .

(cd2) for each index  $k$ ,  $1 \leq k \leq h$ ,  $G_k$  is internally 3-connected.

(cd3) for each  $k$ ,  $2 \leq k \leq h-1$ , vertices in  $V_k$  are on the outer vertices and the following (a) and (b) holds.



## Canonical Decomposition

(cd1)  $V_1$  is the set of all vertices on the inner face containing edge  $(u_1, u_2)$ .

(cd2) for each index  $k$ ,  $1 \leq k \leq h$ ,  $G_k$  is internally 3-connected.

(cd3) for each  $k$ ,  $2 \leq k \leq h-1$ , vertices in  $V_k$  are on the outer vertices and the following (a) and (b) holds.

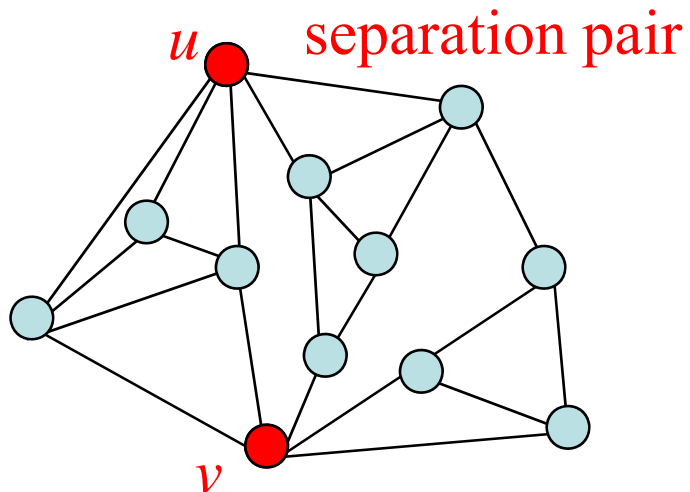
# Internally 3-connected

$G$  is **bi-connected**

For any separation pair  $\{u, v\}$  of  $G$

$u$  and  $v$  are **outer vertices**

each connected component of  $G - \{u, v\}$  contains **an outer vertex**.



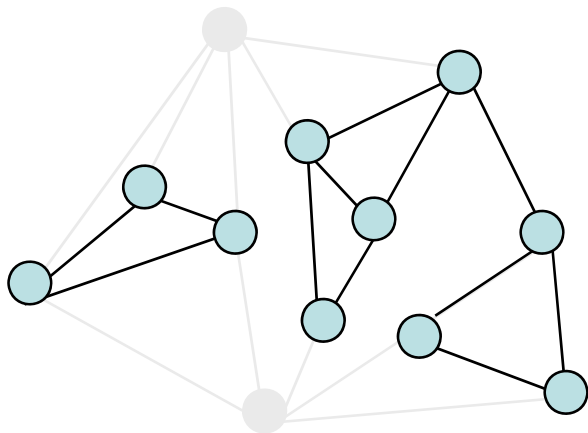
# Internally 3-connected

$G$  is **bi-connected**

For any separation pair  $\{u, v\}$  of  $G$

$u$  and  $v$  are **outer vertices**

each connected component of  $G - \{u, v\}$  contains **an outer vertex**.



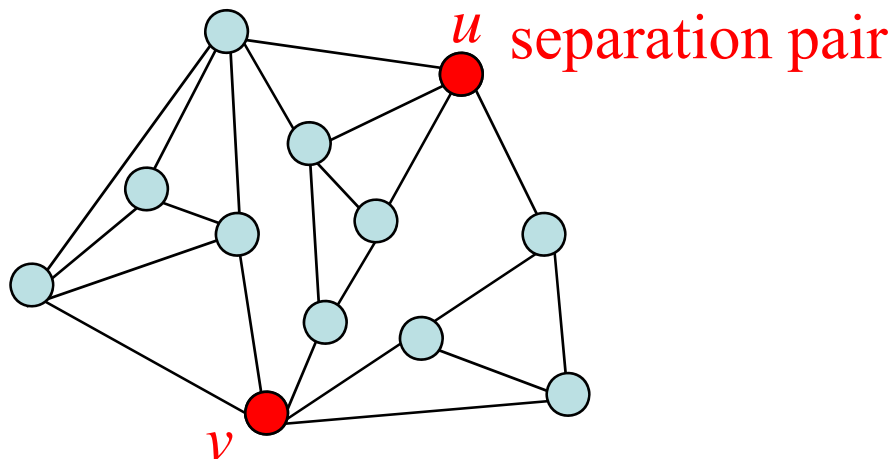
# Internally 3-connected

$G$  is **bi-connected**

For any separation pair  $\{u, v\}$  of  $G$

$u$  and  $v$  are **outer vertices**

each connected component of  $G - \{u, v\}$  contains **an outer vertex**.



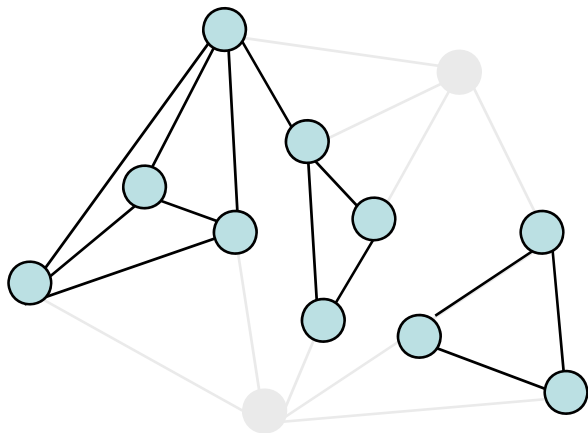
# Internally 3-connected

$G$  is **bi-connected**

For any separation pair  $\{u, v\}$  of  $G$

$u$  and  $v$  are **outer vertices**

each connected component of  $G - \{u, v\}$  contains **an outer vertex**.



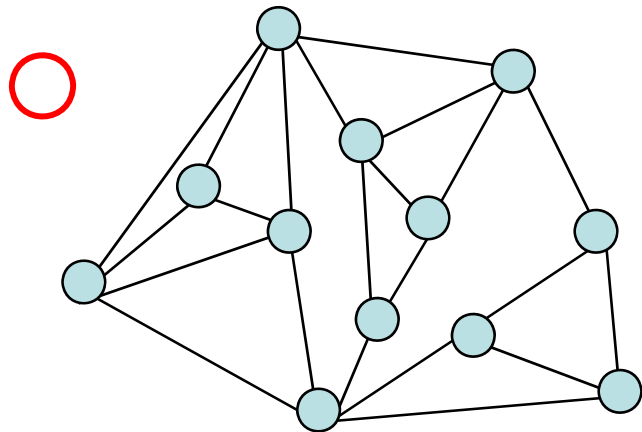
# Internally 3-connected

$G$  is **bi-connected**

For any separation pair  $\{u, v\}$  of  $G$

$u$  and  $v$  are **outer vertices**

each connected component of  $G - \{u, v\}$  contains **an outer vertex**.



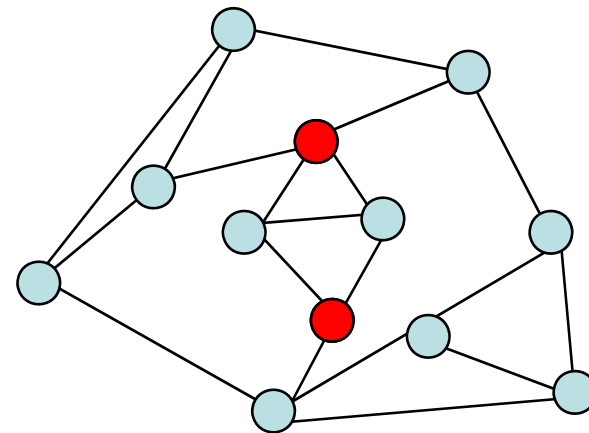
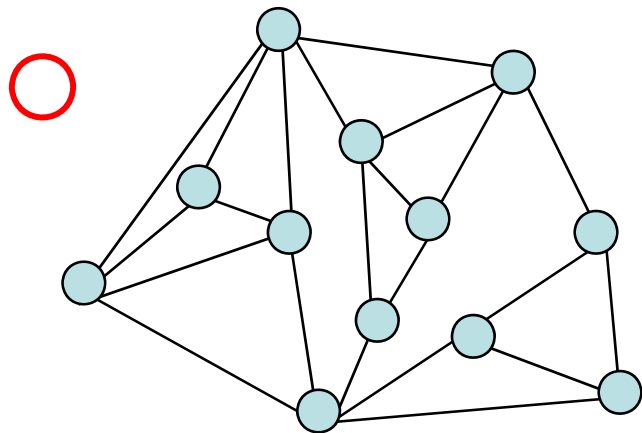
# Internally 3-connected

$G$  is **bi-connected**

For any separation pair  $\{u, v\}$  of  $G$

$u$  and  $v$  are **outer vertices**

each connected component of  $G - \{u, v\}$  contains **an outer vertex**.





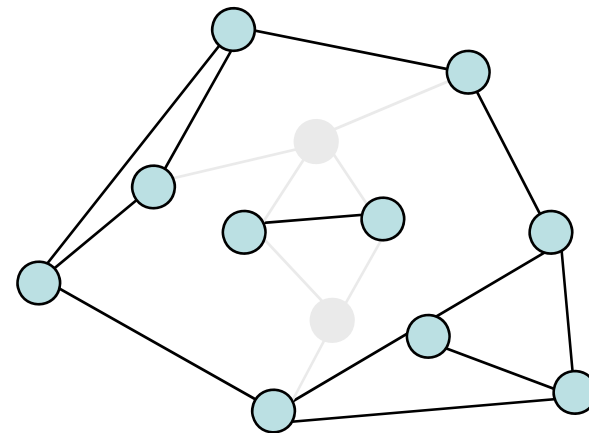
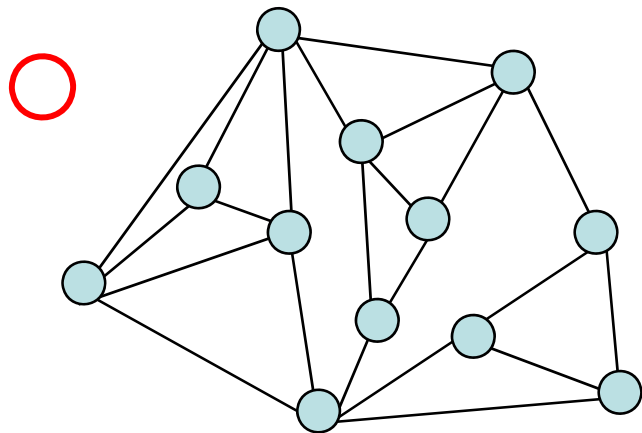
# Internally 3-connected

$G$  is **bi-connected**

For any separation pair  $\{u, v\}$  of  $G$

$u$  and  $v$  are **outer vertices**

each connected component of  $G - \{u, v\}$  contains **an outer vertex**.



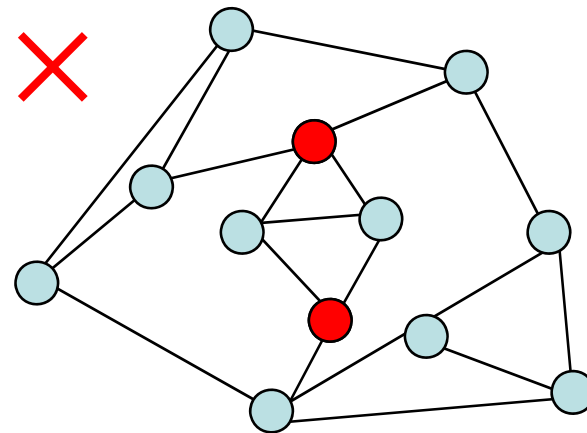
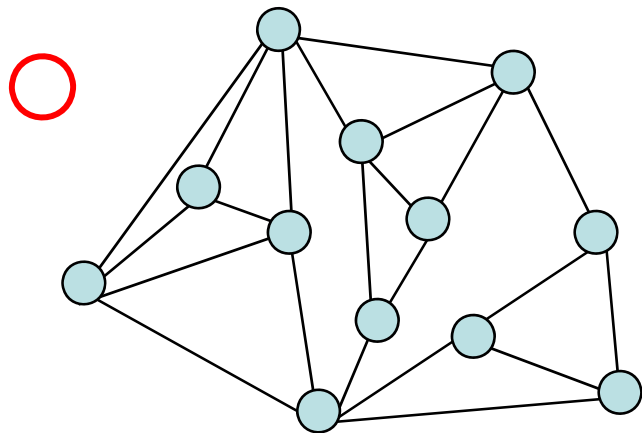
# Internally 3-connected

$G$  is **bi-connected**

For any separation pair  $\{u, v\}$  of  $G$

$u$  and  $v$  are **outer vertices**

each connected component of  $G - \{u, v\}$  contains **an outer vertex**.



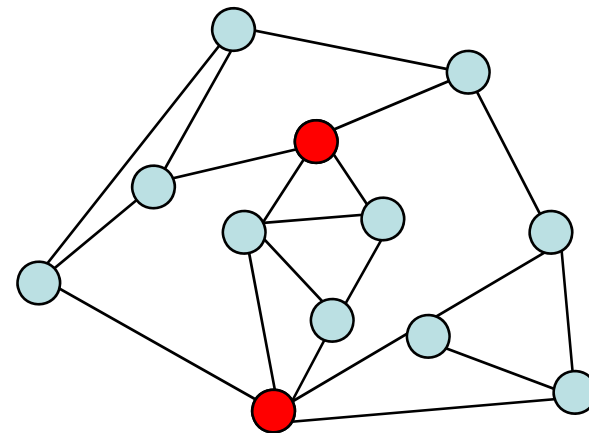
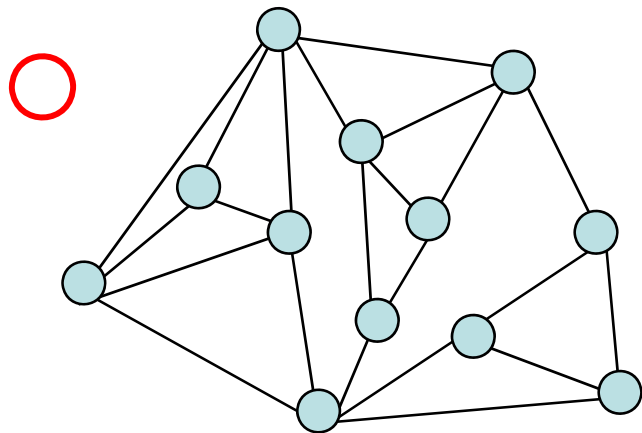
# Internally 3-connected

$G$  is **bi-connected**

For any separation pair  $\{u, v\}$  of  $G$

$u$  and  $v$  are **outer vertices**

each connected component of  $G - \{u, v\}$  contains **an outer vertex**.



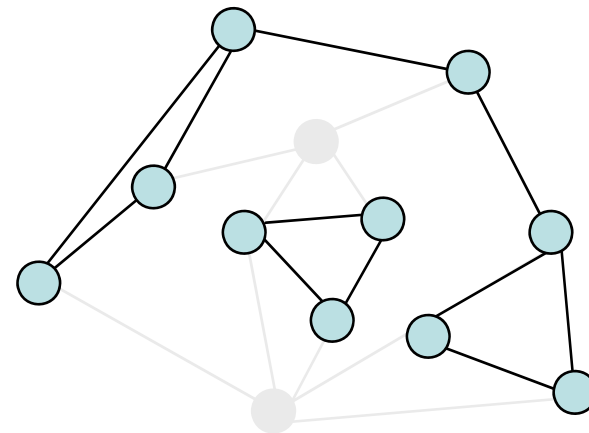
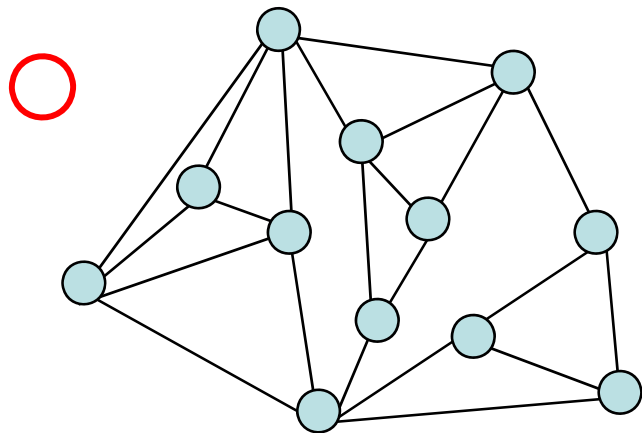
# Internally 3-connected

$G$  is **bi-connected**

For any separation pair  $\{u, v\}$  of  $G$

$u$  and  $v$  are **outer vertices**

each connected component of  $G - \{u, v\}$  contains **an outer vertex**.



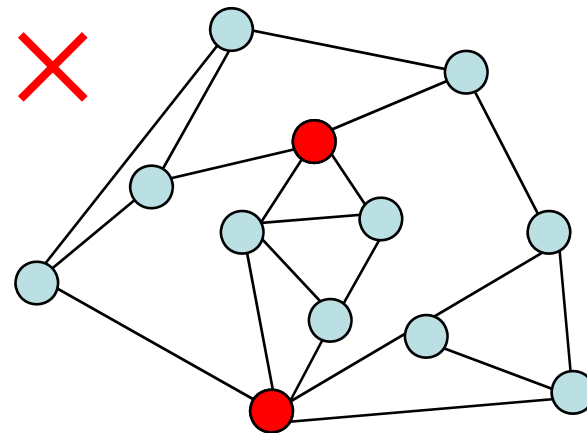
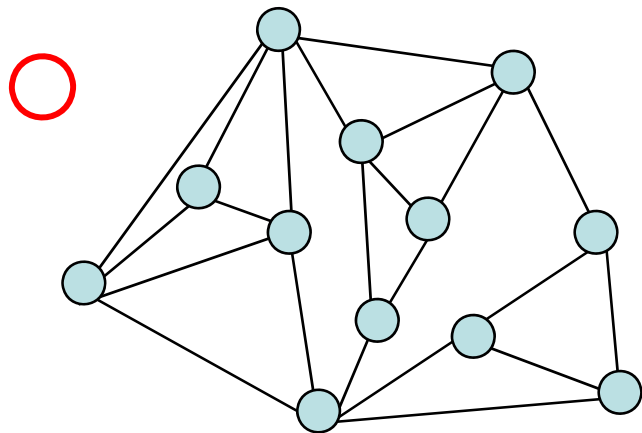
# Internally 3-connected

$G$  is **bi-connected**

For any separation pair  $\{u, v\}$  of  $G$

$u$  and  $v$  are **outer vertices**

each connected component of  $G - \{u, v\}$  contains **an outer vertex**.



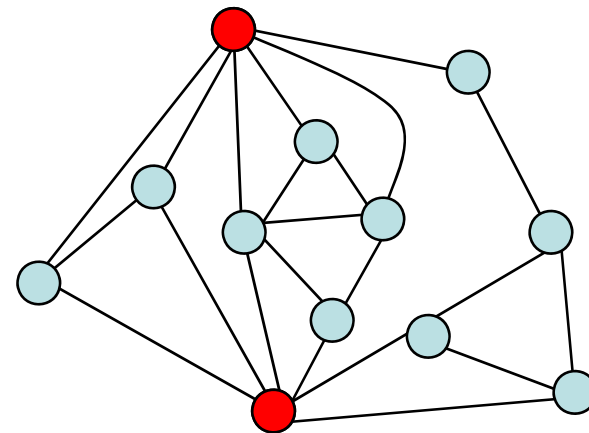
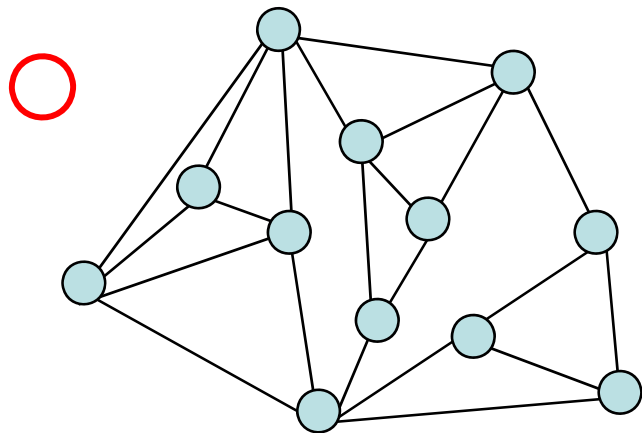
# Internally 3-connected

$G$  is **bi-connected**

For any separation pair  $\{u, v\}$  of  $G$

$u$  and  $v$  are **outer vertices**

each connected component of  $G - \{u, v\}$  contains **an outer vertex**.



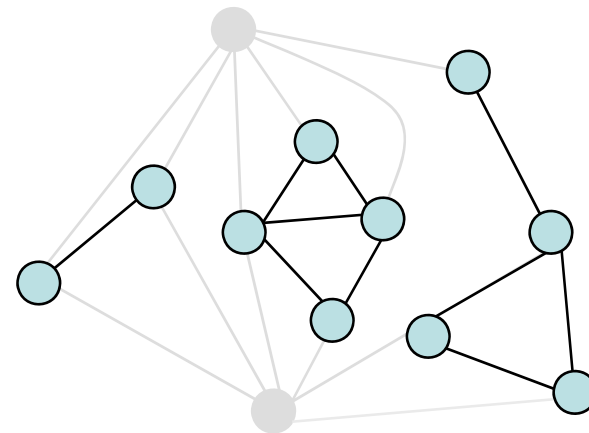
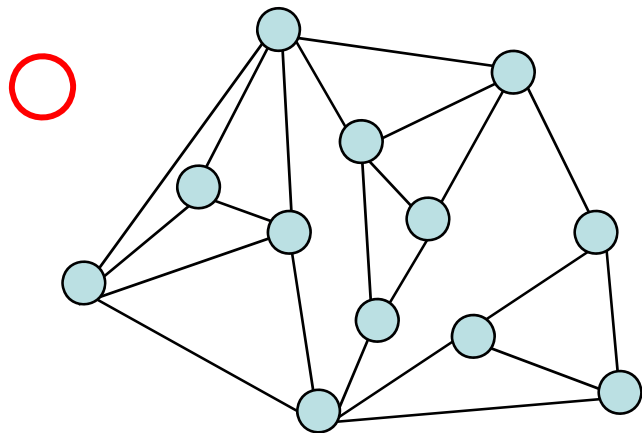
# Internally 3-connected

$G$  is **bi-connected**

For any separation pair  $\{u, v\}$  of  $G$

$u$  and  $v$  are **outer vertices**

each connected component of  $G - \{u, v\}$  contains **an outer vertex**.



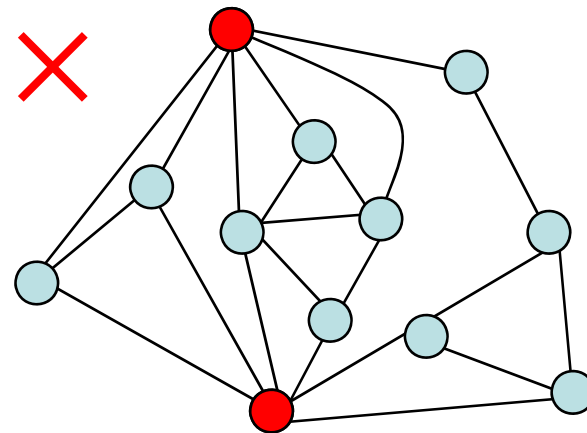
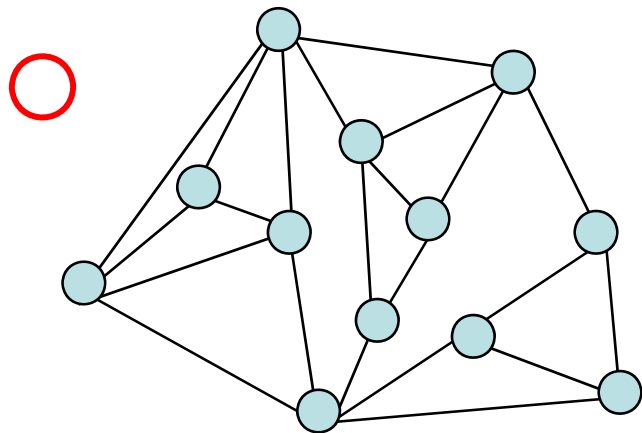
# Internally 3-connected

$G$  is **bi-connected**

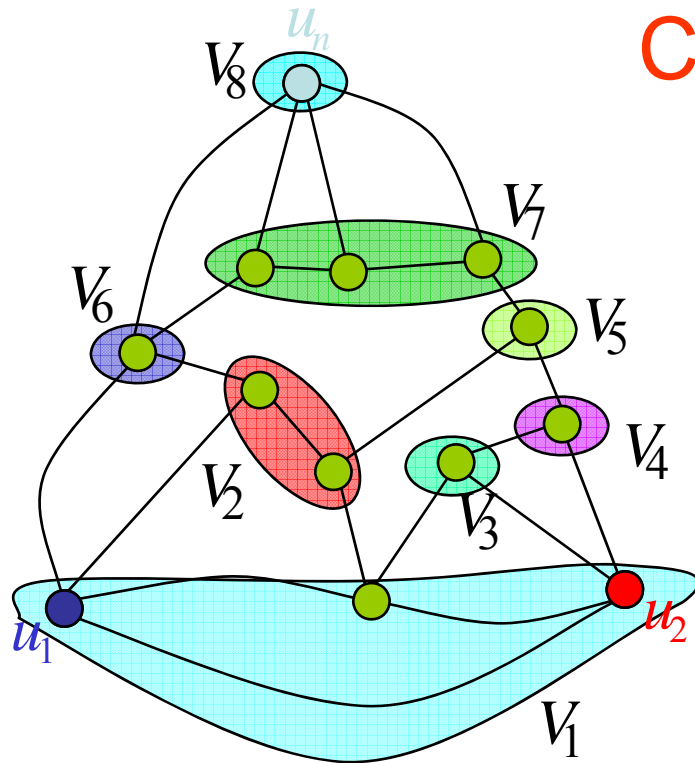
For any separation pair  $\{u, v\}$  of  $G$

$u$  and  $v$  are **outer vertices**

each connected component of  $G - \{u, v\}$  contains **an outer vertex**.





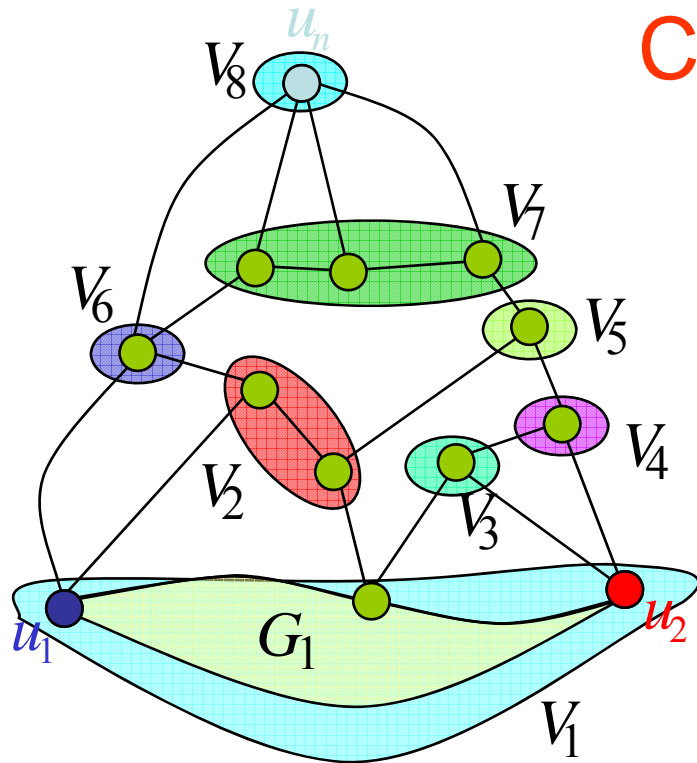


## Canonical Decomposition

(cd1)  $V_1$  is the set of all vertices on the inner face containing edge  $(u_1, u_2)$ .

(cd2) for each index  $k$ ,  $1 \leq k \leq h$ ,  $G_k$  is internally 3-connected.

(cd3) for each  $k$ ,  $2 \leq k \leq h-1$ , vertices in  $V_k$  are on the outer vertices and the following (a) and (b) holds.

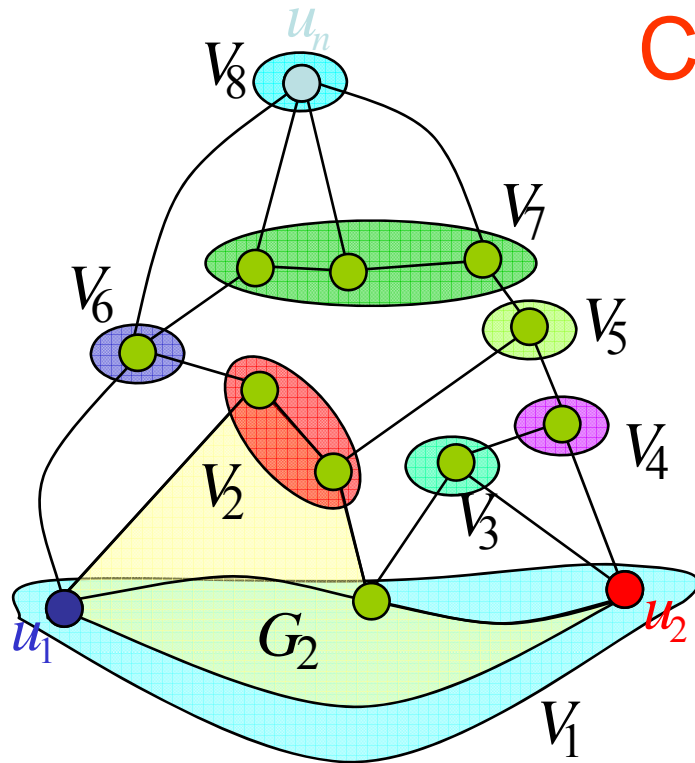


## Canonical Decomposition

(cd1)  $V_1$  is the set of all vertices on the inner face containing edge  $(u_1, u_2)$ .

(cd2) for each index  $k$ ,  $1 \leq k \leq h$ ,  $G_k$  is internally 3-connected.

(cd3) for each  $k$ ,  $2 \leq k \leq h-1$ , vertices in  $V_k$  are on the outer vertices and the following (a) and (b) holds.

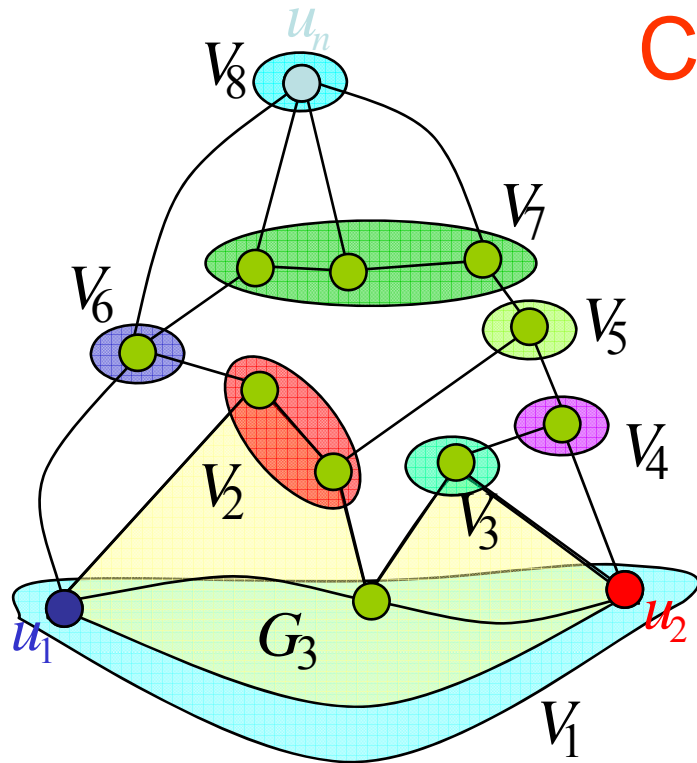


## Canonical Decomposition

(cd1)  $V_1$  is the set of all vertices on the inner face containing edge  $(u_1, u_2)$ .

(cd2) for each index  $k$ ,  $1 \leq k \leq h$ ,  $G_k$  is internally 3-connected.

(cd3) for each  $k$ ,  $2 \leq k \leq h-1$ , vertices in  $V_k$  are on the outer vertices and the following (a) and (b) holds.

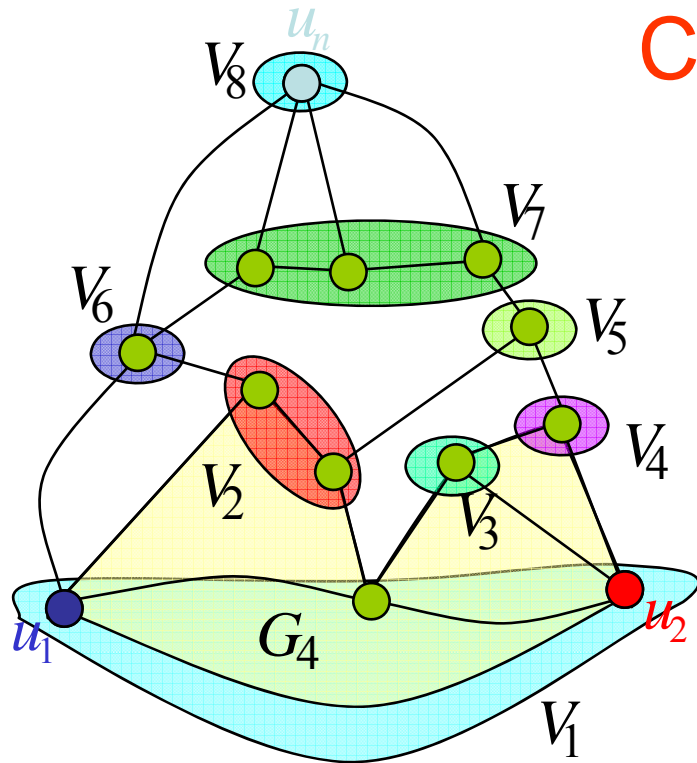


## Canonical Decomposition

(cd1)  $V_1$  is the set of all vertices on the inner face containing edge  $(u_1, u_2)$ .

(cd2) for each index  $k$ ,  $1 \leq k \leq h$ ,  $G_k$  is internally 3-connected.

(cd3) for each  $k$ ,  $2 \leq k \leq h-1$ , vertices in  $V_k$  are on the outer vertices and the following (a) and (b) holds.

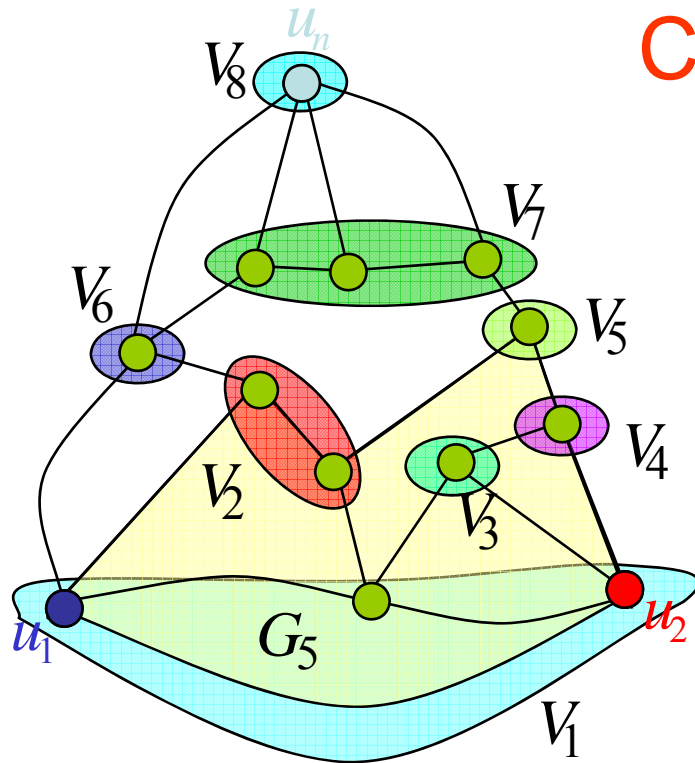


## Canonical Decomposition

(cd1)  $V_1$  is the set of all vertices on the inner face containing edge  $(u_1, u_2)$ .

(cd2) for each index  $k$ ,  $1 \leq k \leq h$ ,  $G_k$  is internally 3-connected.

(cd3) for each  $k$ ,  $2 \leq k \leq h-1$ , vertices in  $V_k$  are on the outer vertices and the following (a) and (b) holds.

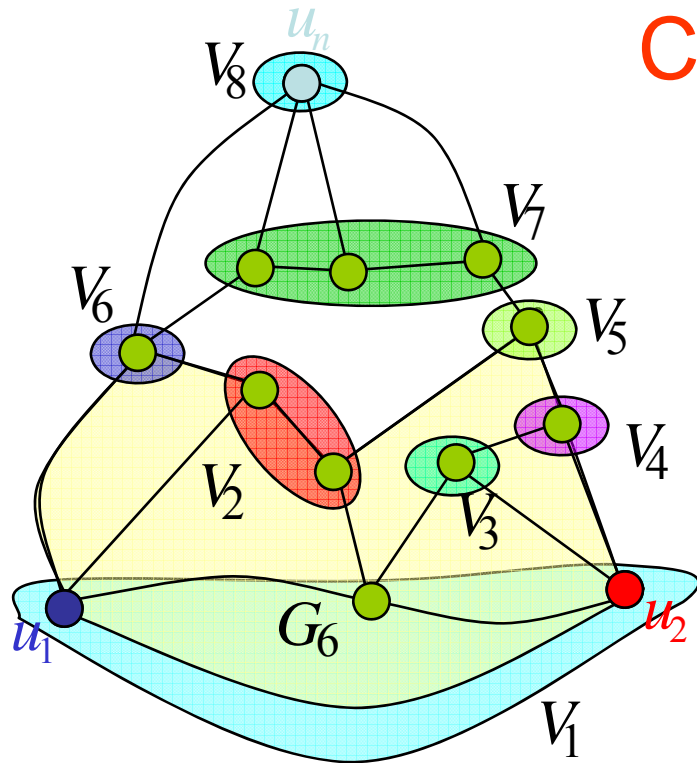


## Canonical Decomposition

(cd1)  $V_1$  is the set of all vertices on the inner face containing edge  $(u_1, u_2)$ .

(cd2) for each index  $k$ ,  $1 \leq k \leq h$ ,  $G_k$  is internally 3-connected.

(cd3) for each  $k$ ,  $2 \leq k \leq h-1$ , vertices in  $V_k$  are on the outer vertices and the following (a) and (b) holds.

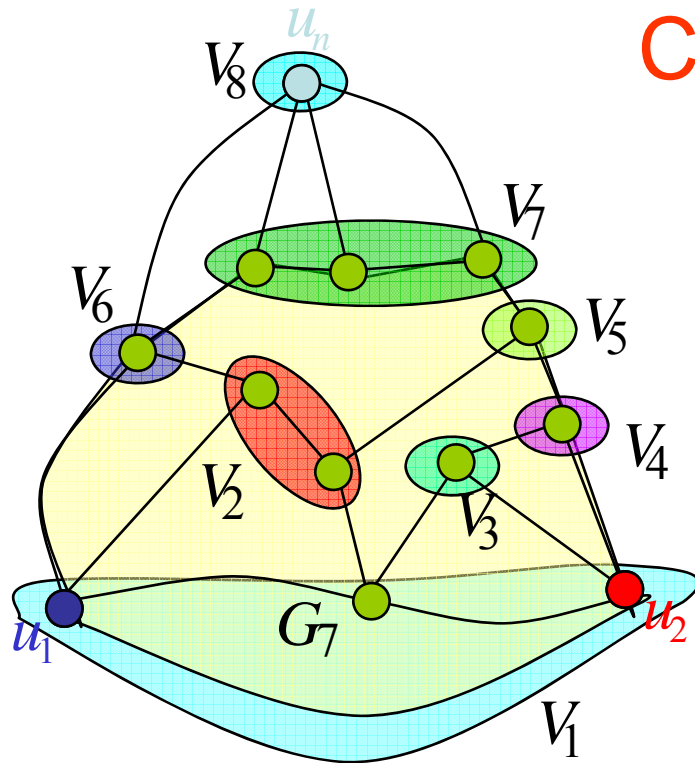


## Canonical Decomposition

(cd1)  $V_1$  is the set of all vertices on the inner face containing edge  $(u_1, u_2)$ .

(cd2) for each index  $k$ ,  $1 \leq k \leq h$ ,  $G_k$  is internally 3-connected.

(cd3) for each  $k$ ,  $2 \leq k \leq h-1$ , vertices in  $V_k$  are on the outer vertices and the following (a) and (b) holds.



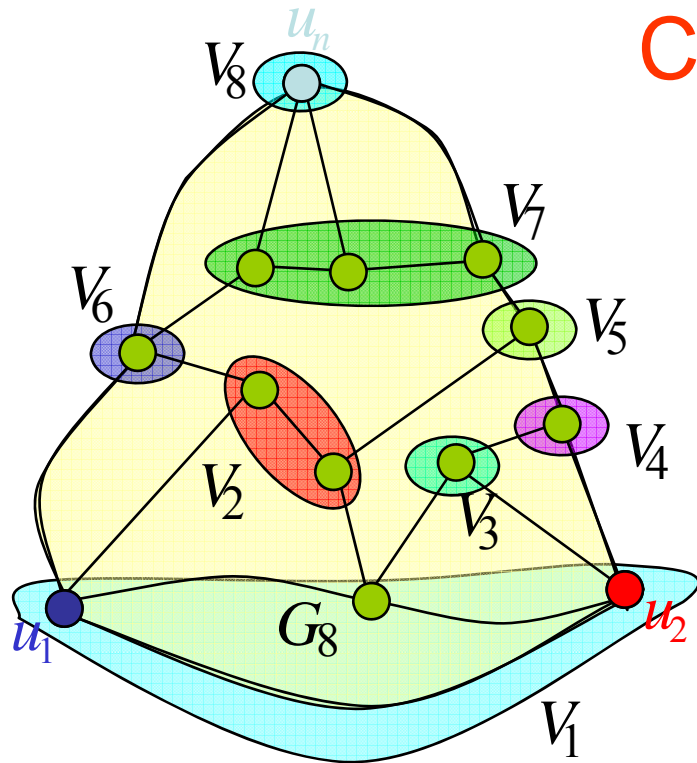
## Canonical Decomposition

(cd1)  $V_1$  is the set of all vertices on the inner face containing edge  $(u_1, u_2)$ .

(cd2) for each index  $k$ ,  $1 \leq k \leq h$ ,  $G_k$  is internally 3-connected.

(cd3) for each  $k$ ,  $2 \leq k \leq h-1$ , vertices in  $V_k$  are on the outer vertices and the following (a) and (b) holds.





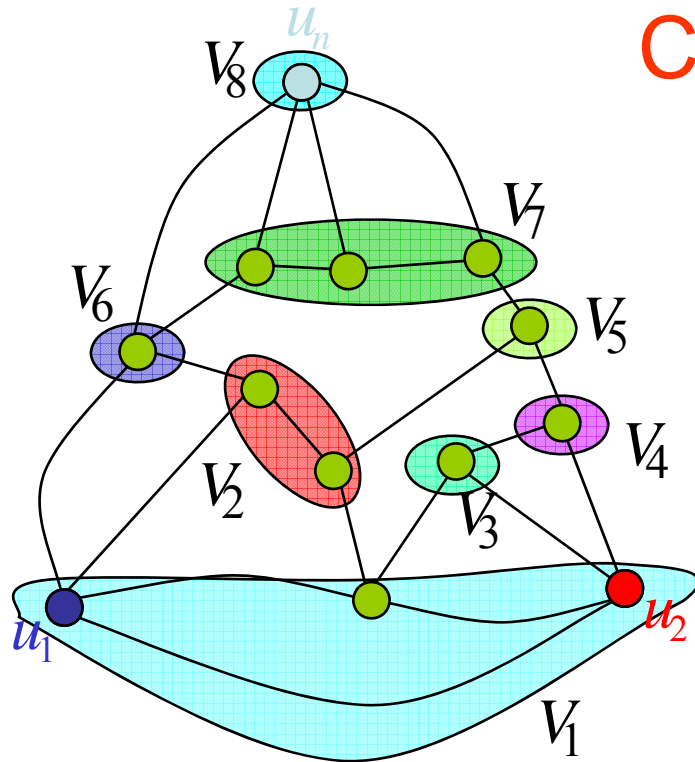
## Canonical Decomposition

(cd1)  $V_1$  is the set of all vertices on the inner face containing edge  $(u_1, u_2)$ .

(cd2) for each index  $k$ ,  $1 \leq k \leq h$ ,  $G_k$  is internally 3-connected.

(cd3) for each  $k$ ,  $2 \leq k \leq h-1$ , vertices in  $V_k$  are on the outer vertices and the following (a) and (b) holds.

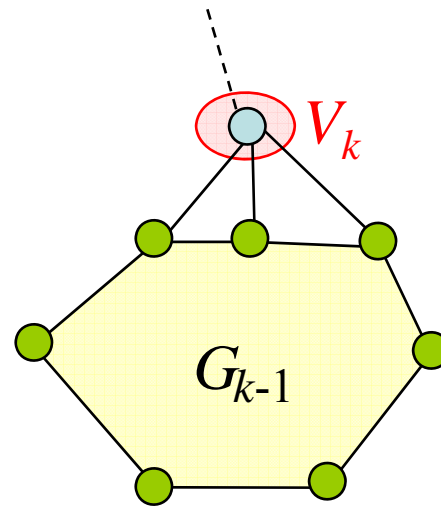
# Canonical Decomposition



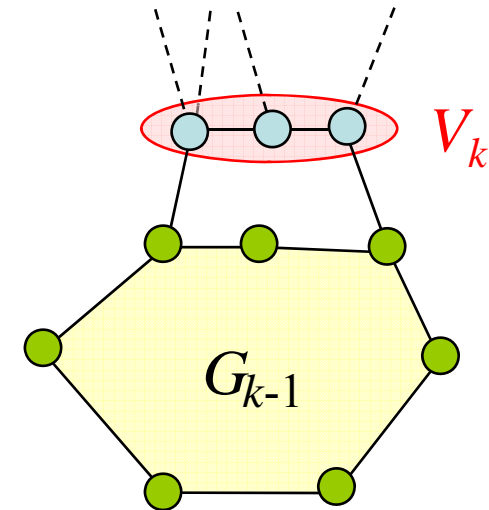
(cd1)  $V_1$  is the set of all vertices on the inner face containing edge  $(u_1, u_2)$ .

(cd2) for each index  $k$ ,  $1 \leq k \leq h$ ,  $G_k$  is internally 3-connected.

(cd3) for each  $k$ ,  $2 \leq k \leq h-1$ , vertices in  $V_k$  are on the outer vertices and the following (a) and (b) holds.

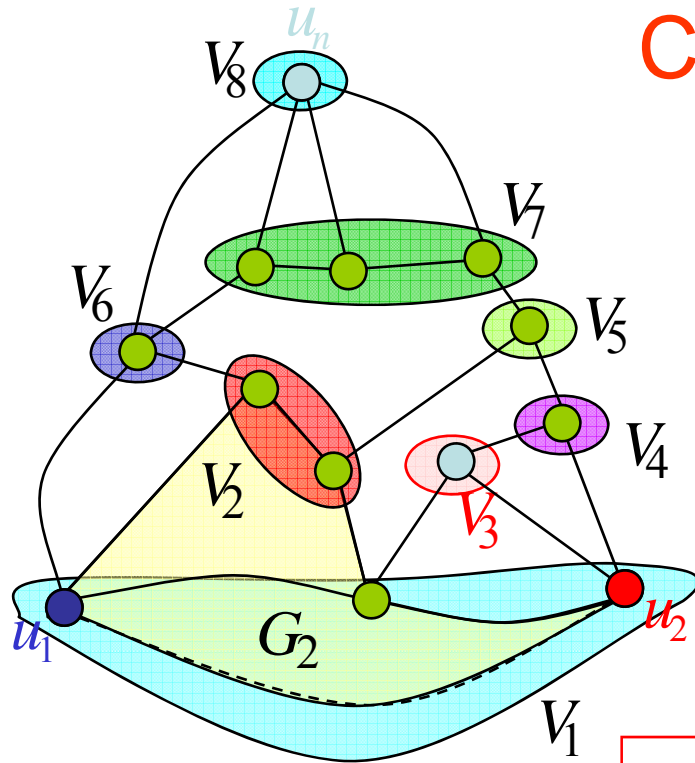


(a)



(b)

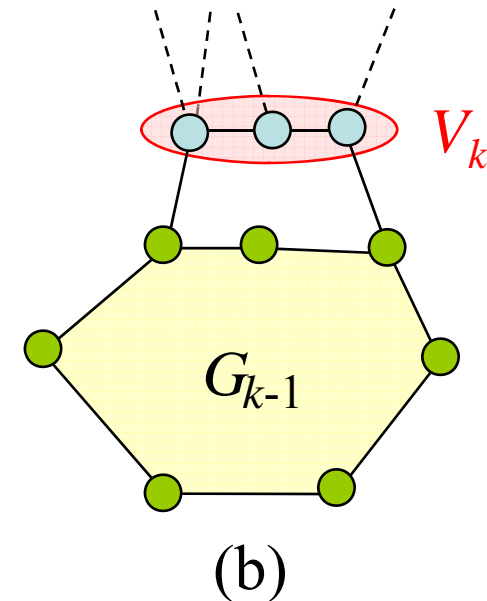
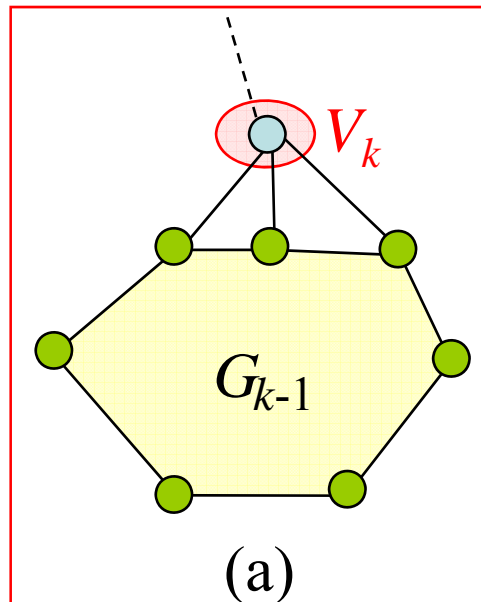
# Canonical Decomposition



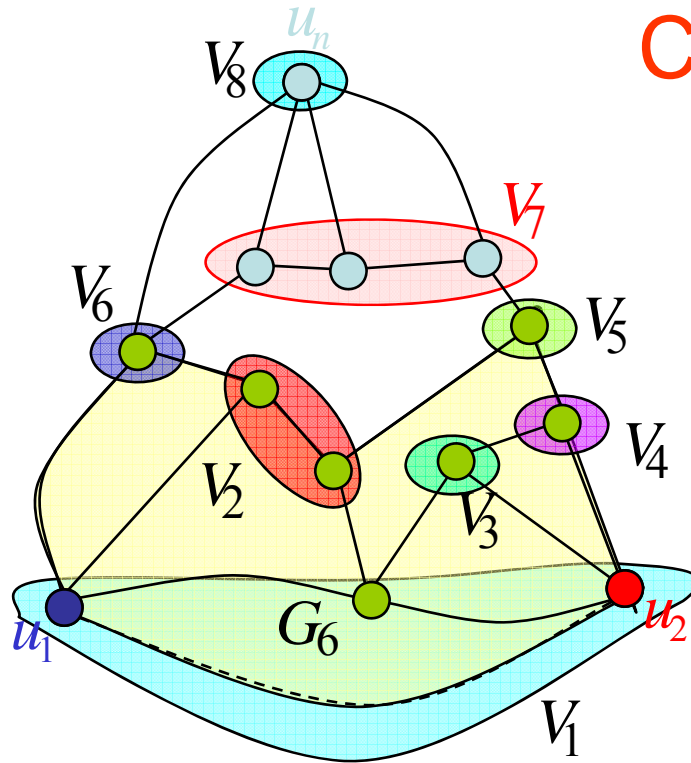
(cd1)  $V_1$  is the set of all vertices on the inner face containing edge  $(u_1, u_2)$ .

(cd2) for each index  $k$ ,  $1 \leq k \leq h$ ,  $G_k$  is internally 3-connected.

(cd3) for each  $k$ ,  $2 \leq k \leq h-1$ , vertices in  $V_k$  are on the outer vertices and the following (a) and (b) holds.



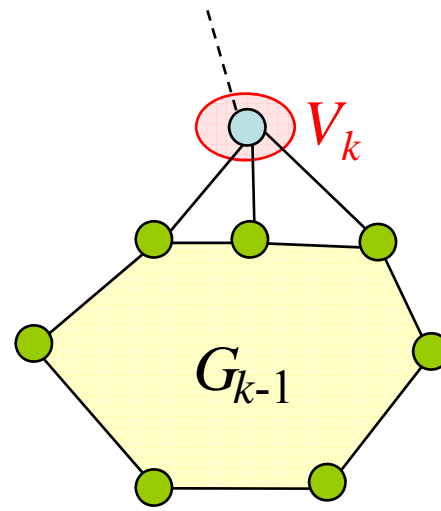
# Canonical Decomposition



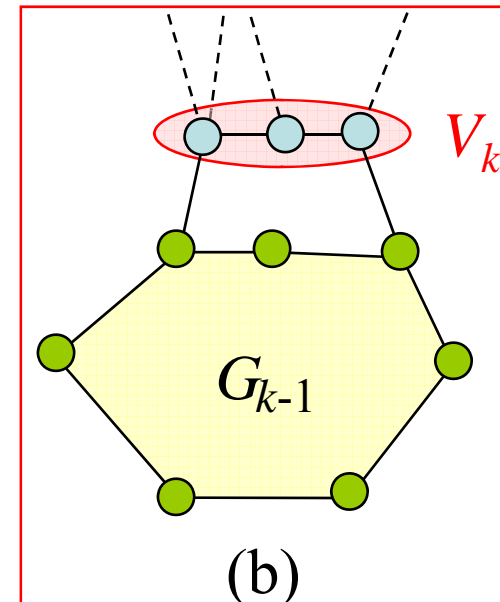
(cd1)  $V_1$  is the set of all vertices on the inner face containing edge  $(u_1, u_2)$ .

(cd2) for each index  $k$ ,  $1 \leq k \leq h$ ,  $G_k$  is internally 3-connected.

(cd3) for each  $k$ ,  $2 \leq k \leq h-1$ , vertices in  $V_k$  are on the outer vertices and the following (a) and (b) holds.



(a)



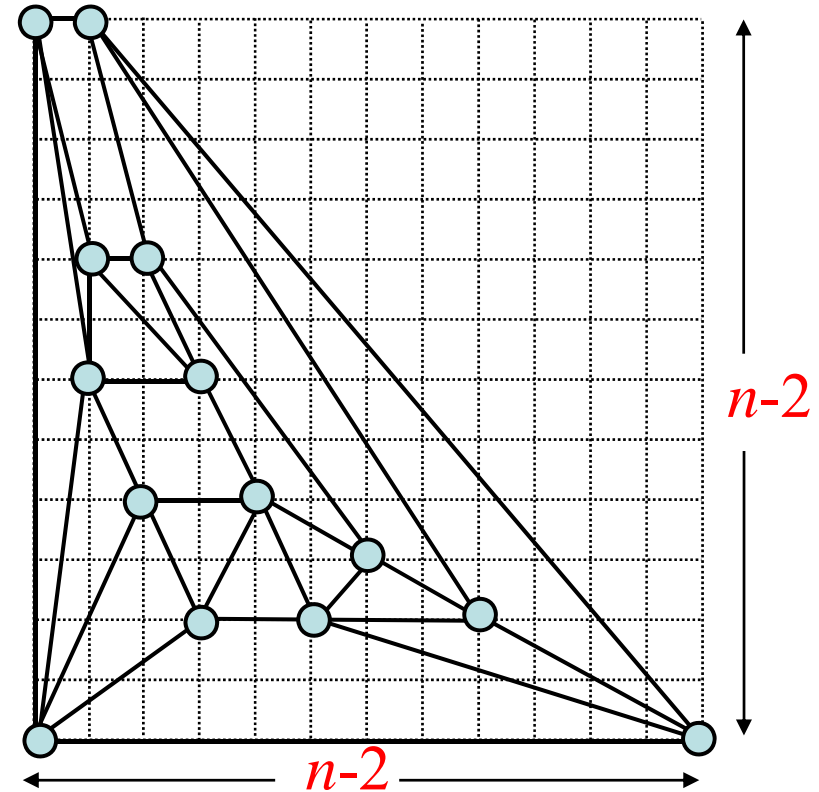
(b)

## Convex Grid Drawing

Chrobak and Kant '97

Input: 3-connected graph

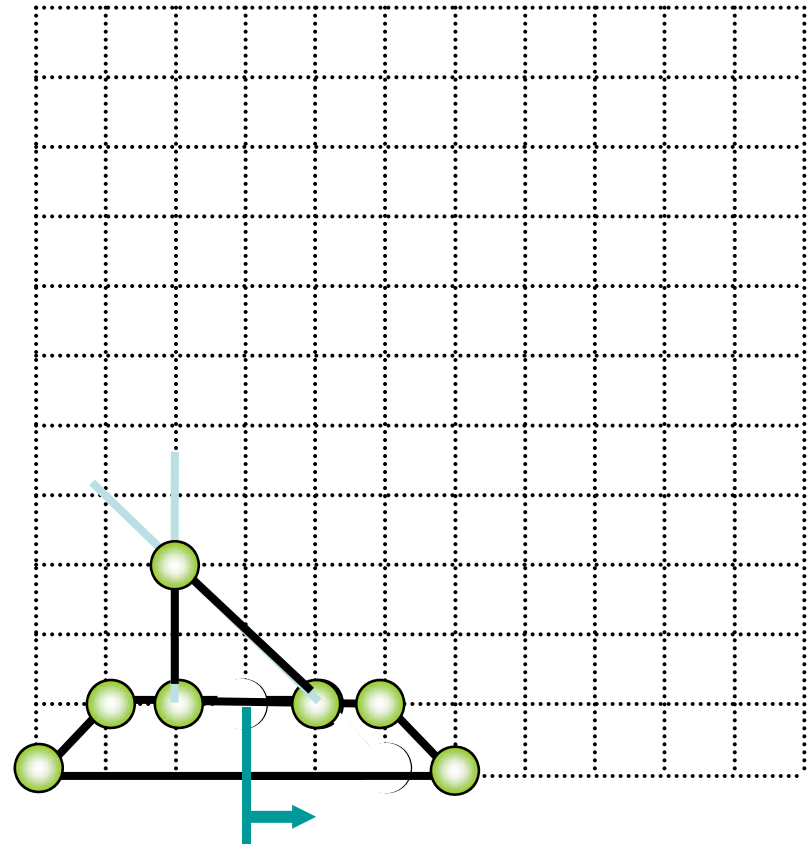
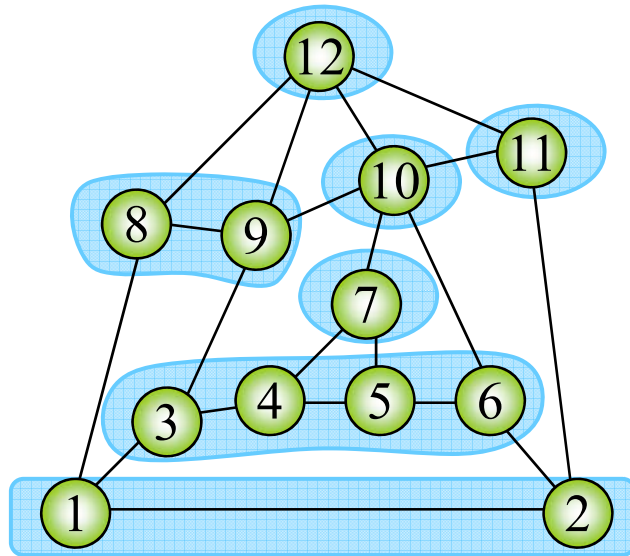
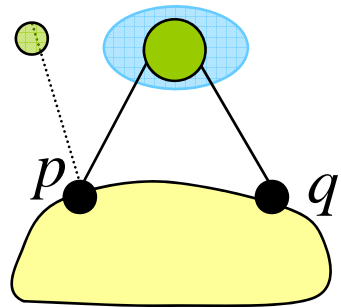
Output: convex grid drawing



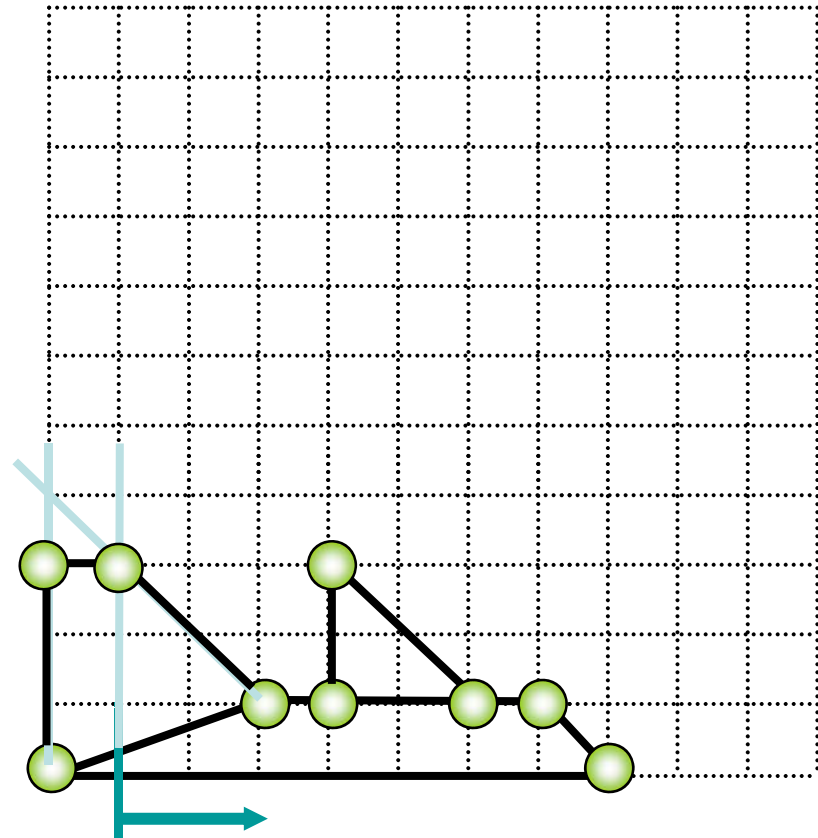
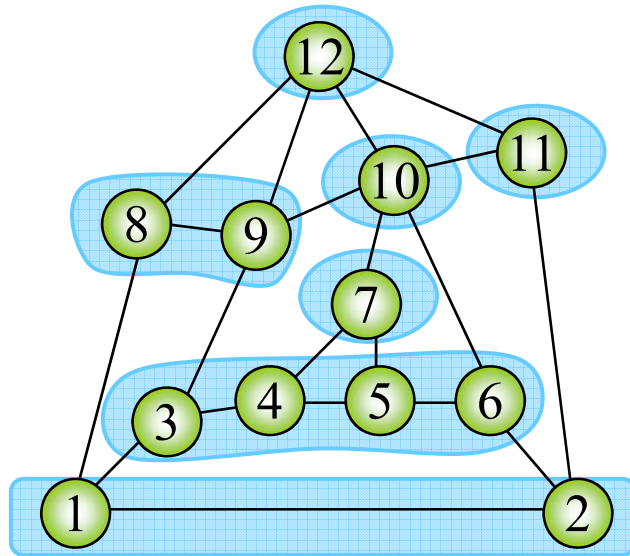
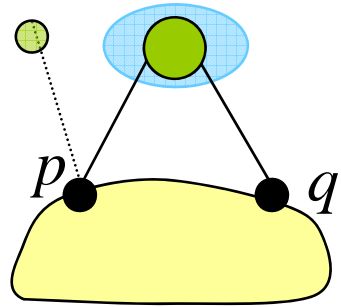
Grid Size

Area  $W \times H \leq n^2$

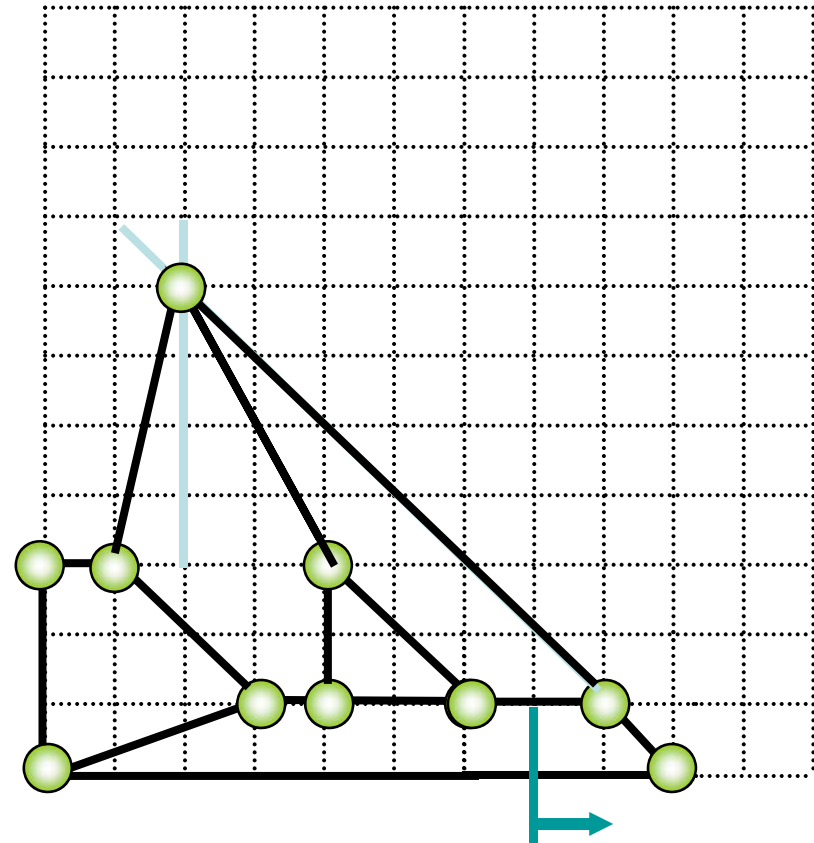
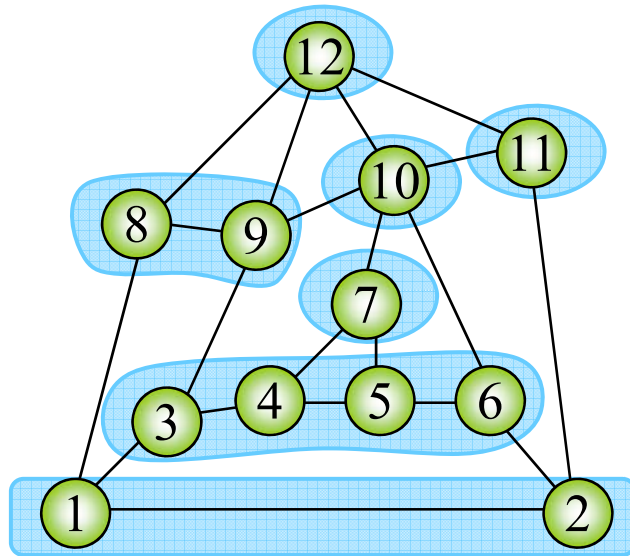
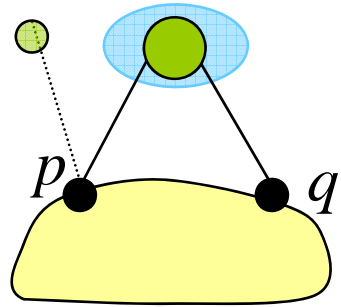
# Shift method



# Shift method

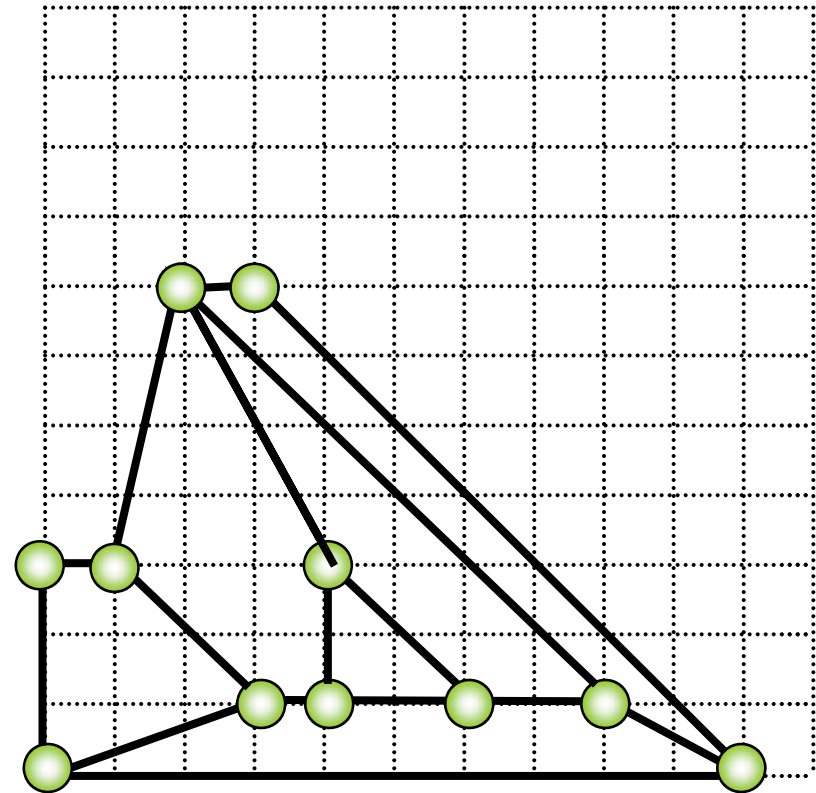
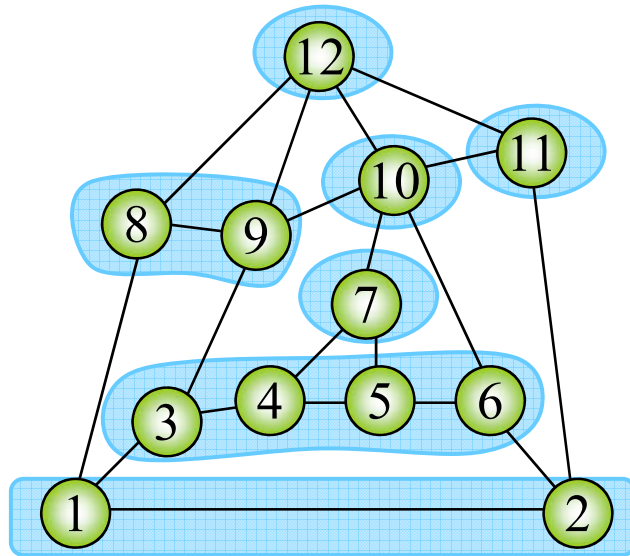
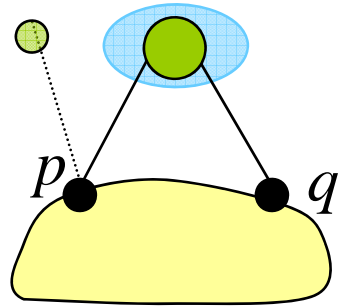


# Shift method

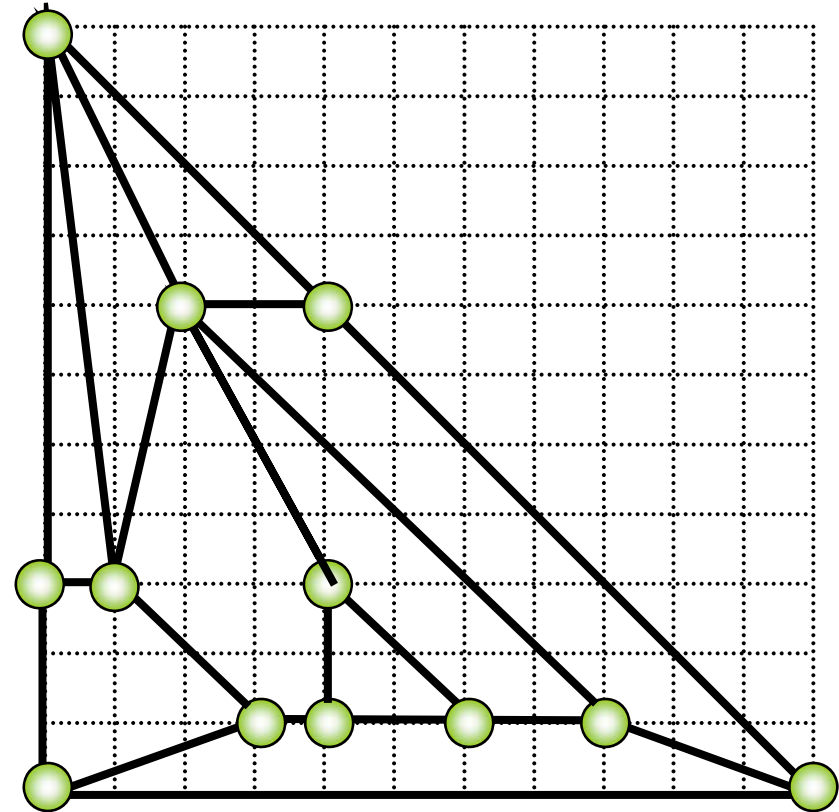
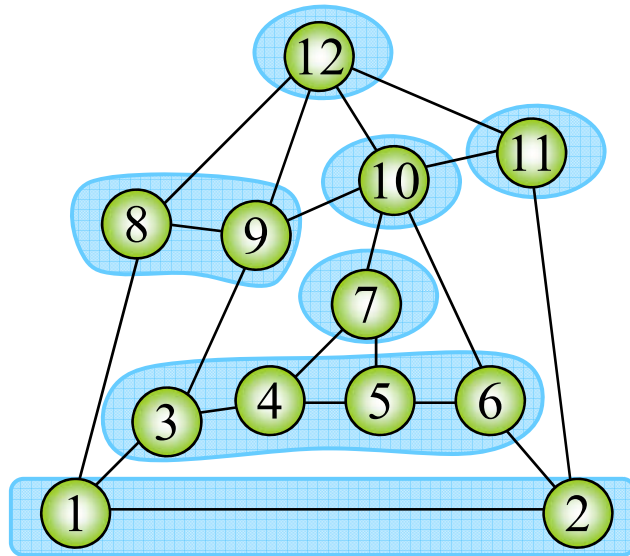
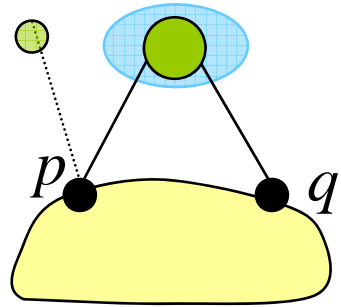




# Shift method

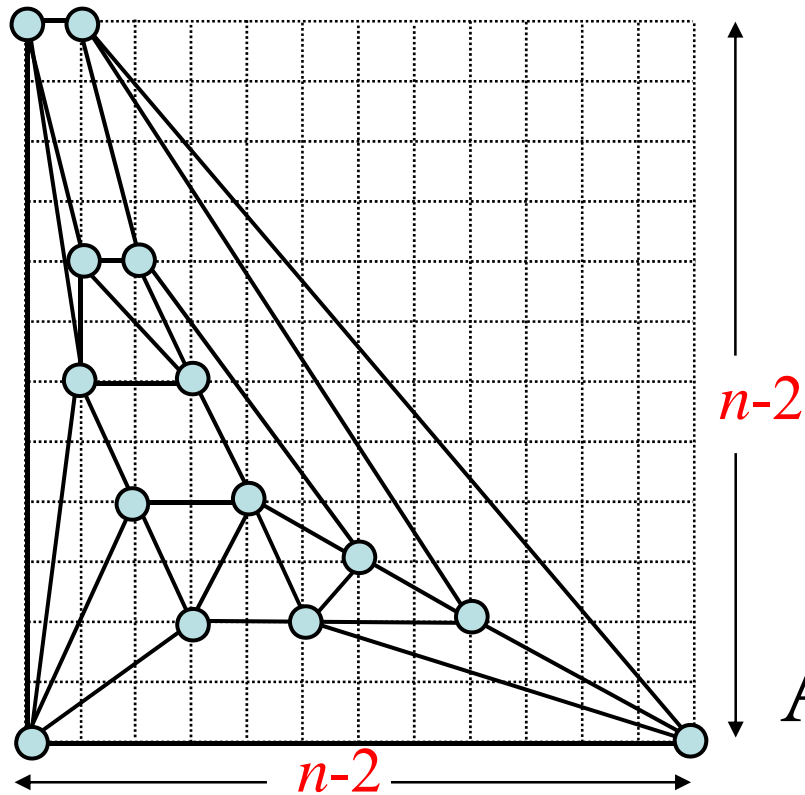


# Shift method



Chrobak and Kant '97

3-connected graph



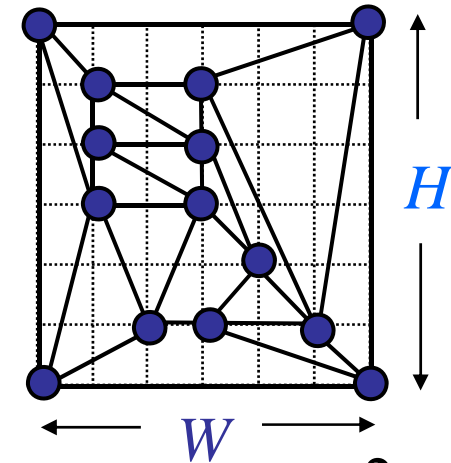
$$\text{Area} \approx n^2$$

Miura *et al.* 2000

4-connected graph

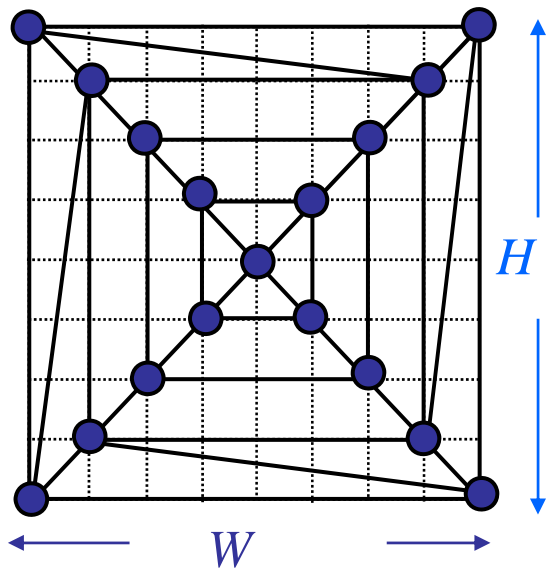


$$\text{Area} \leq 1/4$$



$$\text{Area} \leq \frac{n^2}{4}$$

The algorithm of Miura *et al.* is  
**best possible**

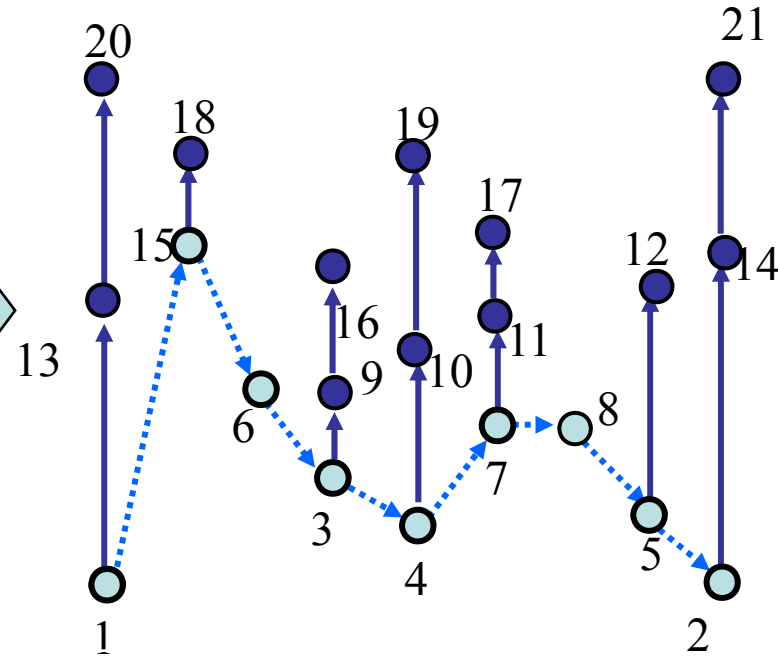
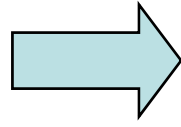
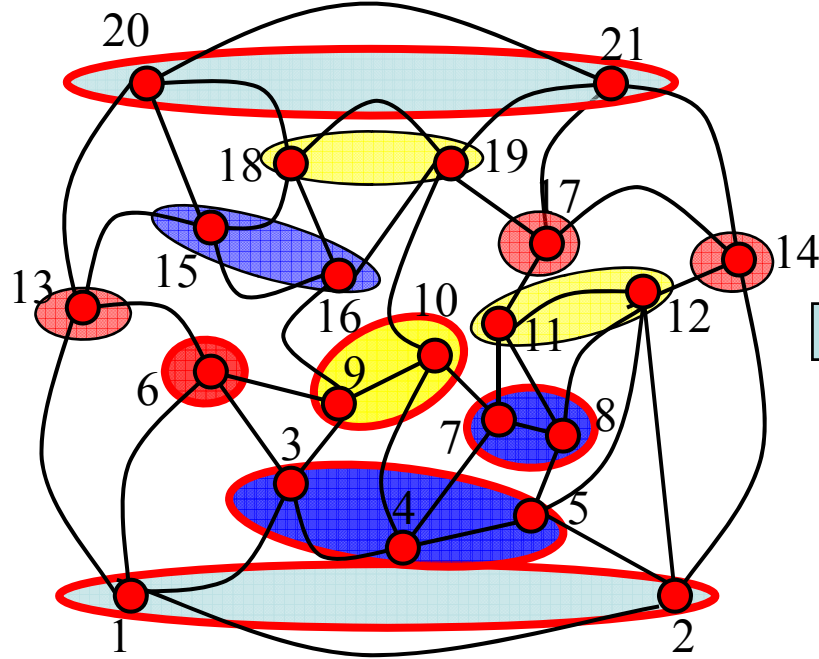


$$W \times H \geq \frac{n^2}{4}$$

## Vertex ordering

- *st*-numbering
- Canonical ordering
- 4-canonical ordering
- Canonical decomposition
- 4-canonical decomposition

# Main idea

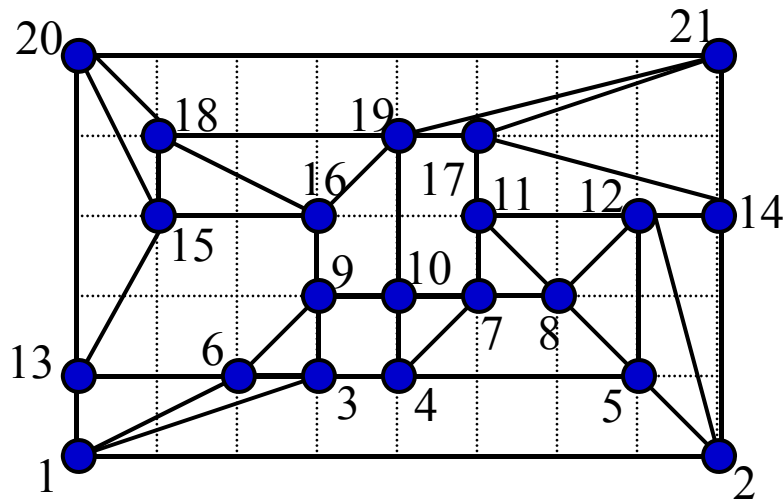


1: 4-canonical decomposition  
 $O(n)$ [NRN97]

2: Find paths



3: Decide x-coordinates

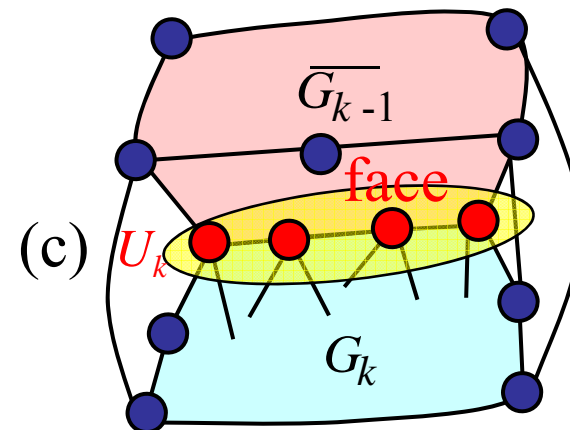
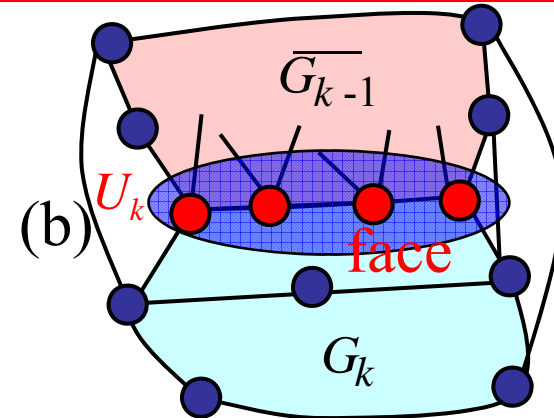
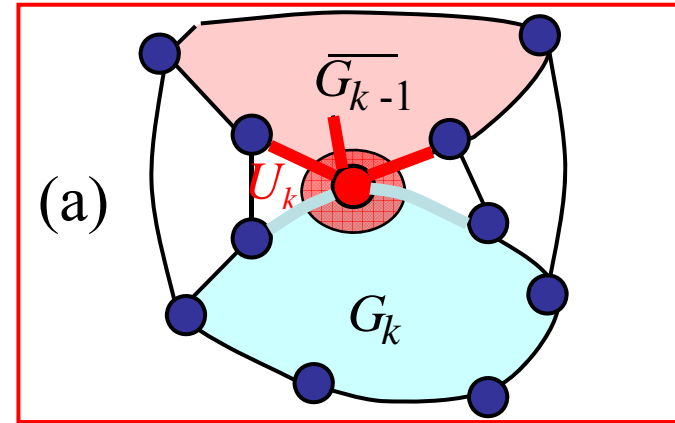
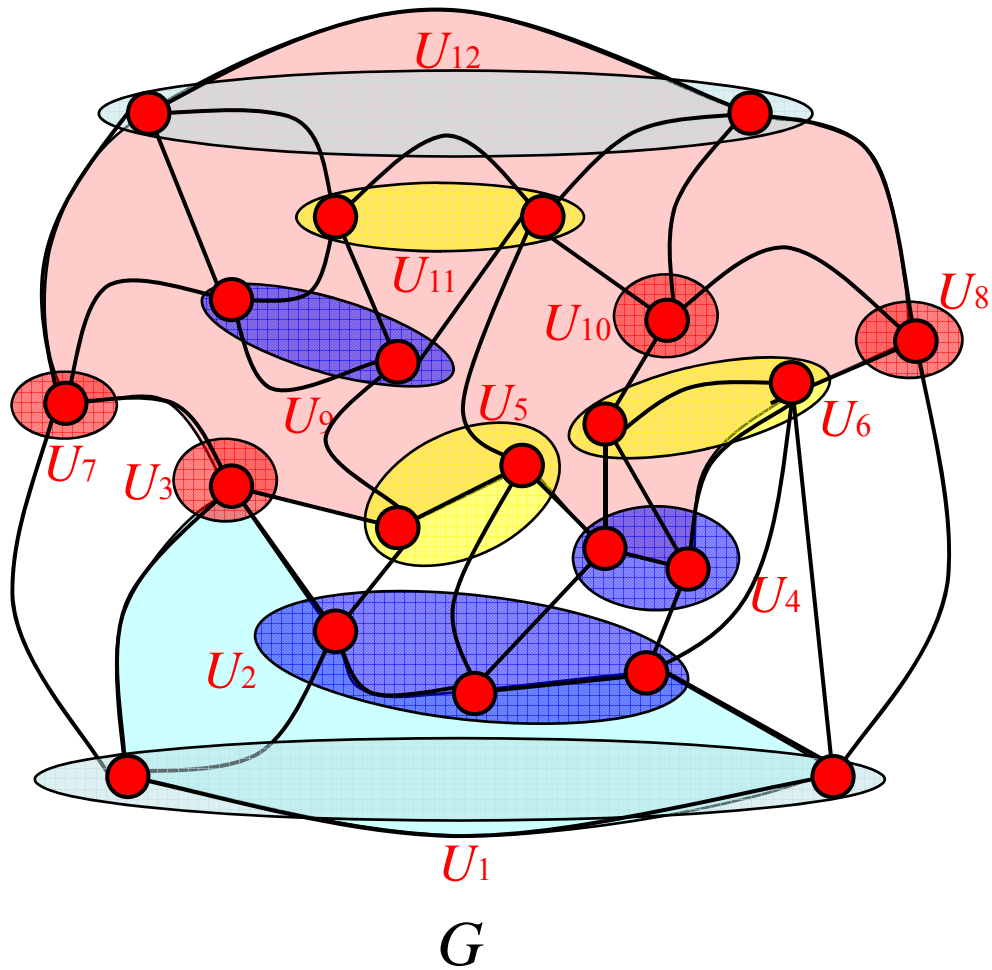


4: Decide y-coordinates

Time complexity:  $O(n)$

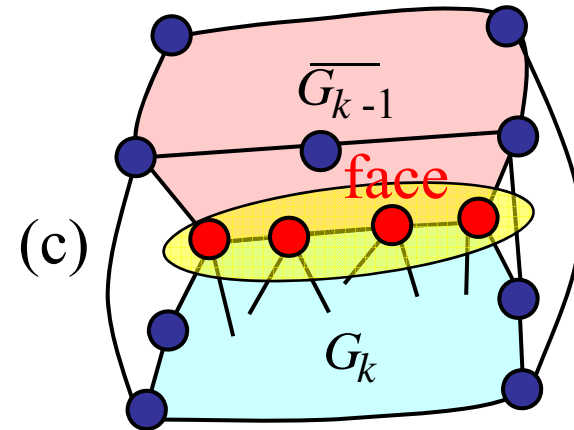
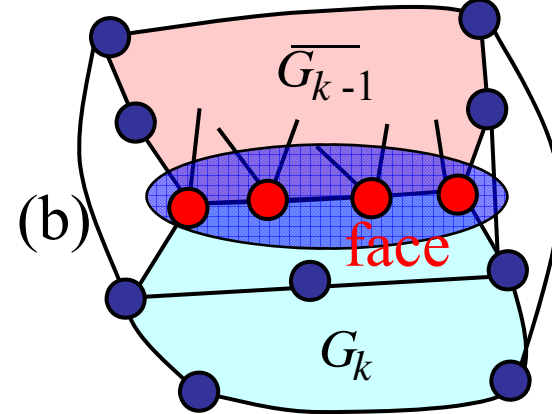
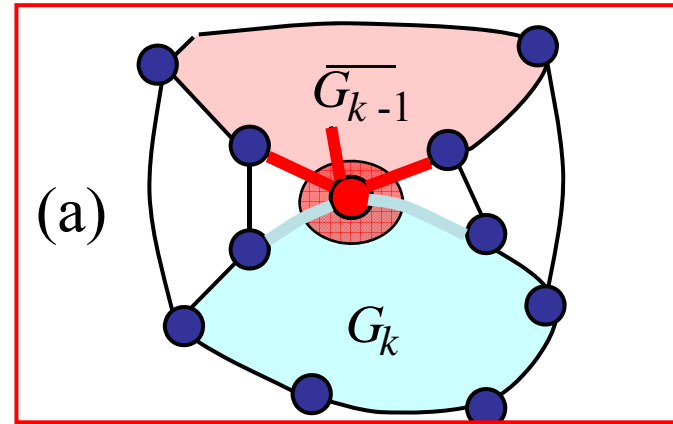
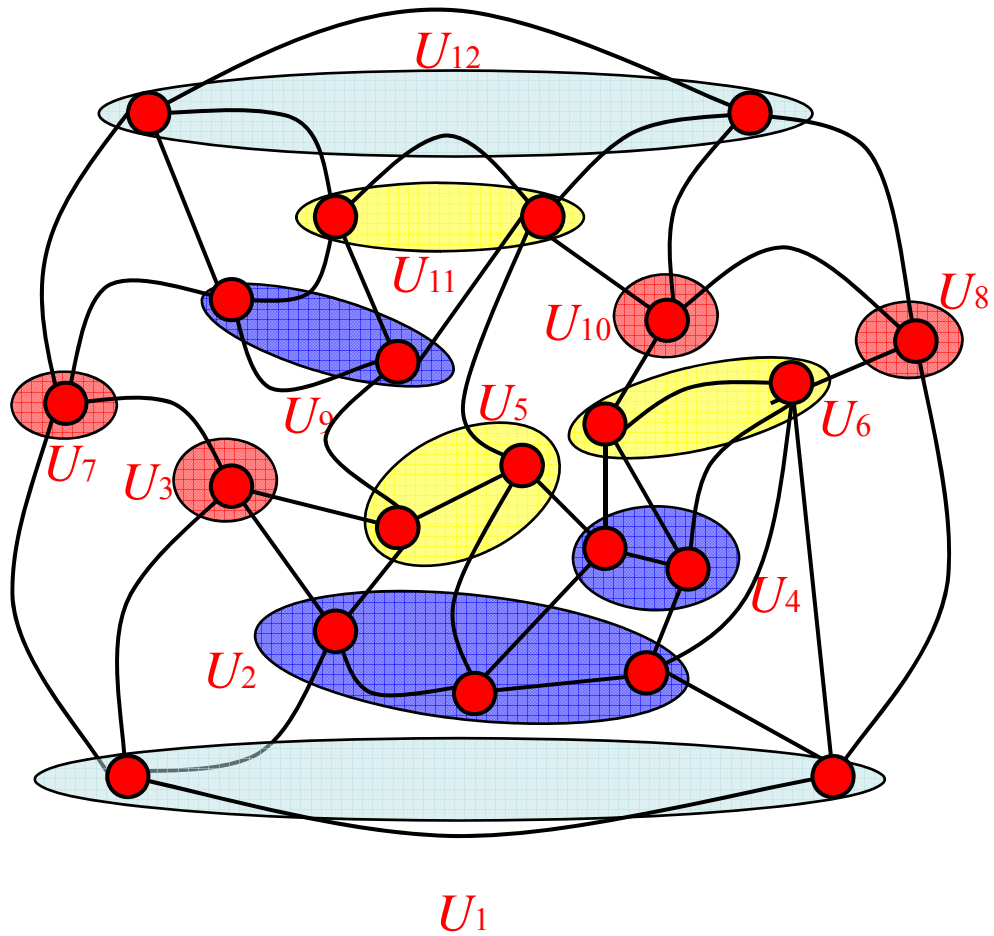
# 4-canonical decomposition [NRN97]

(a generalization of *st*-numbering)



# 4-canonical decomposition [NRN97]

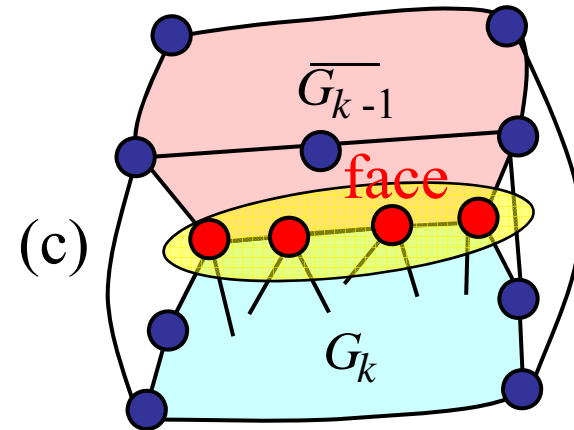
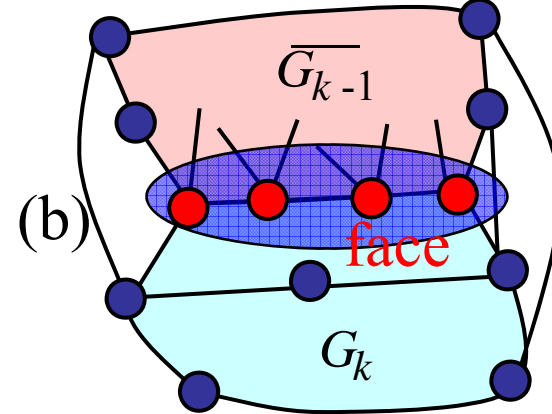
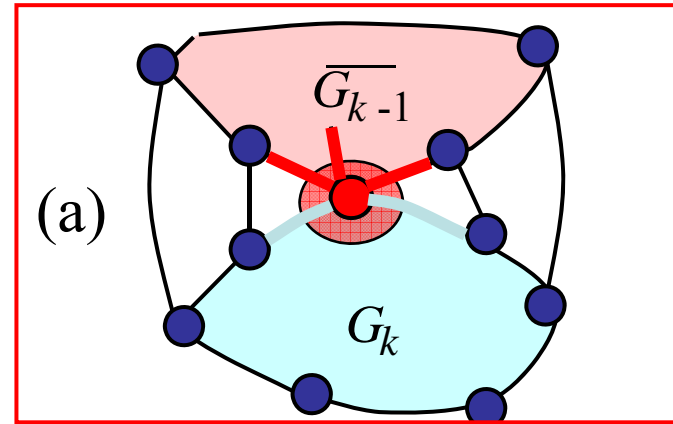
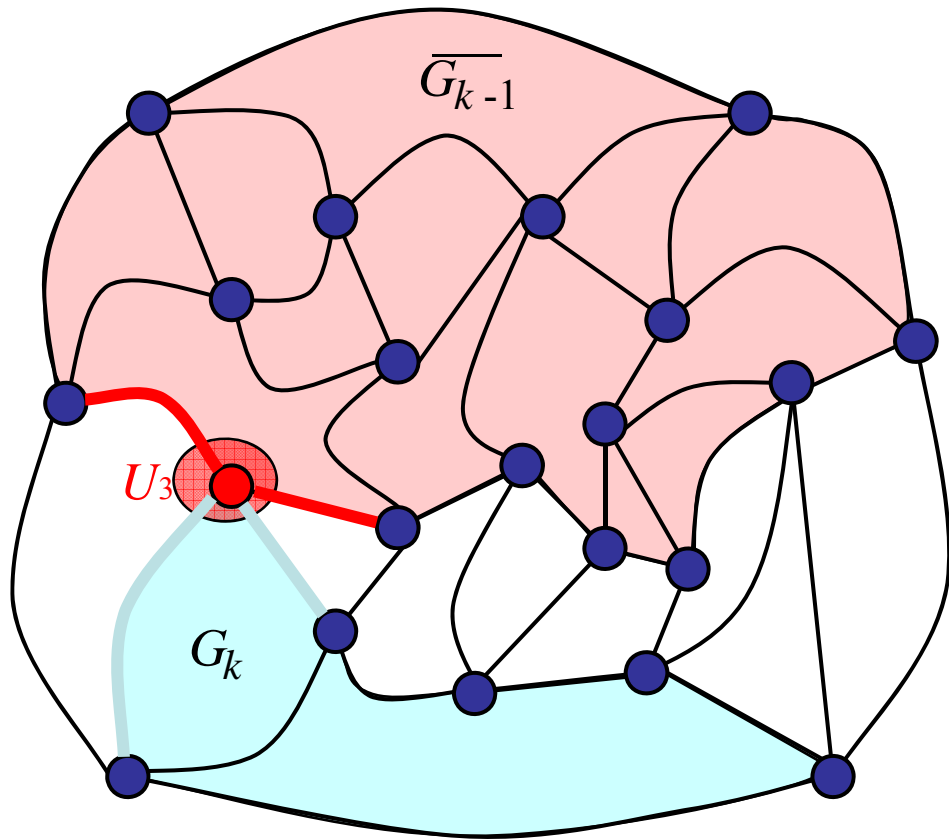
(a generalization of *st*-numbering)





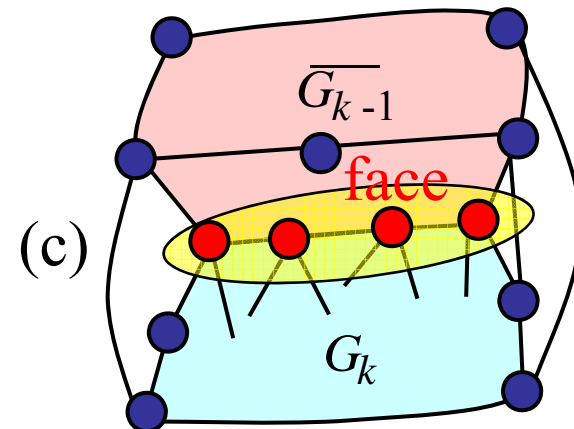
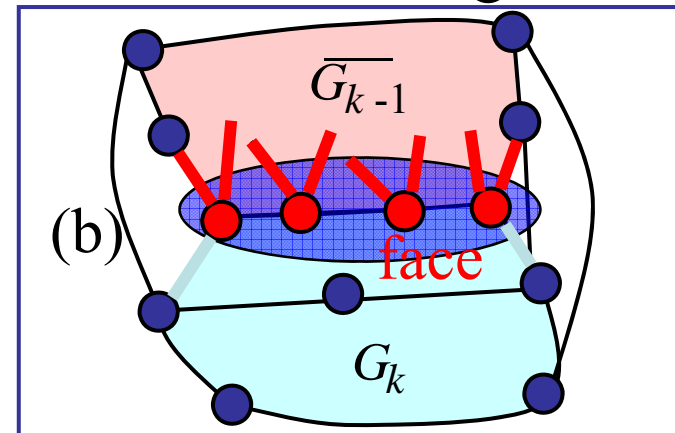
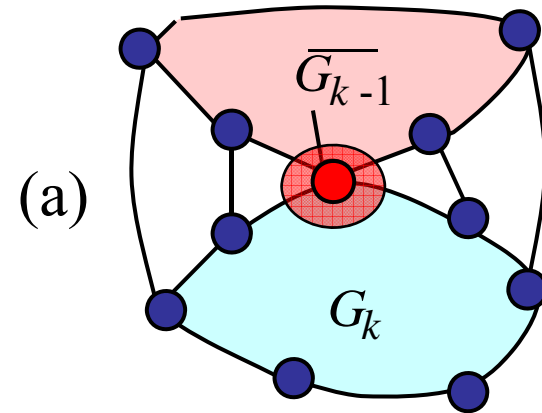
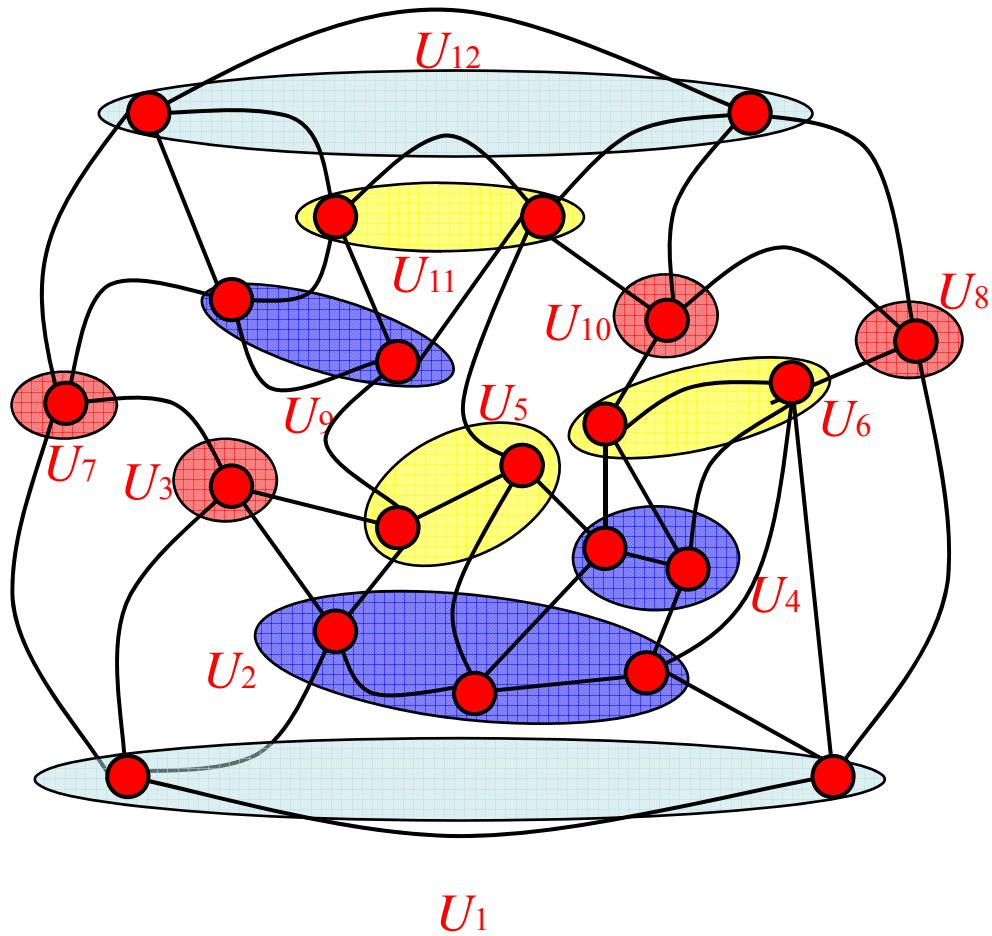
# 4-canonical decomposition [NRN97]

(a generalization of *st*-numbering)



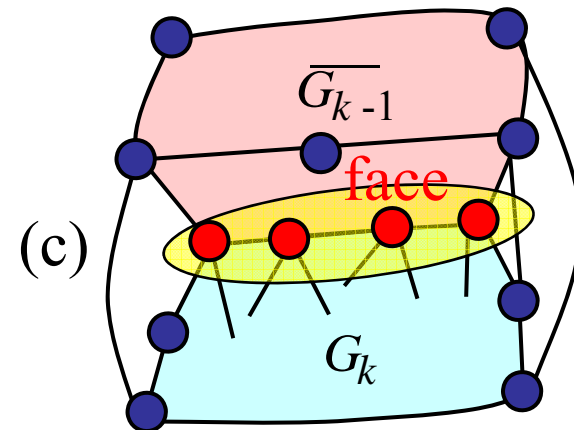
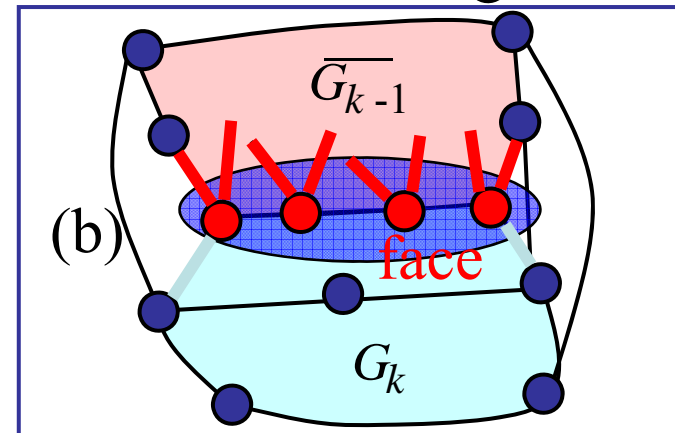
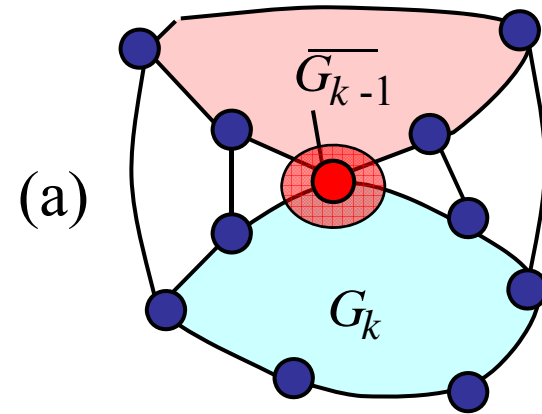
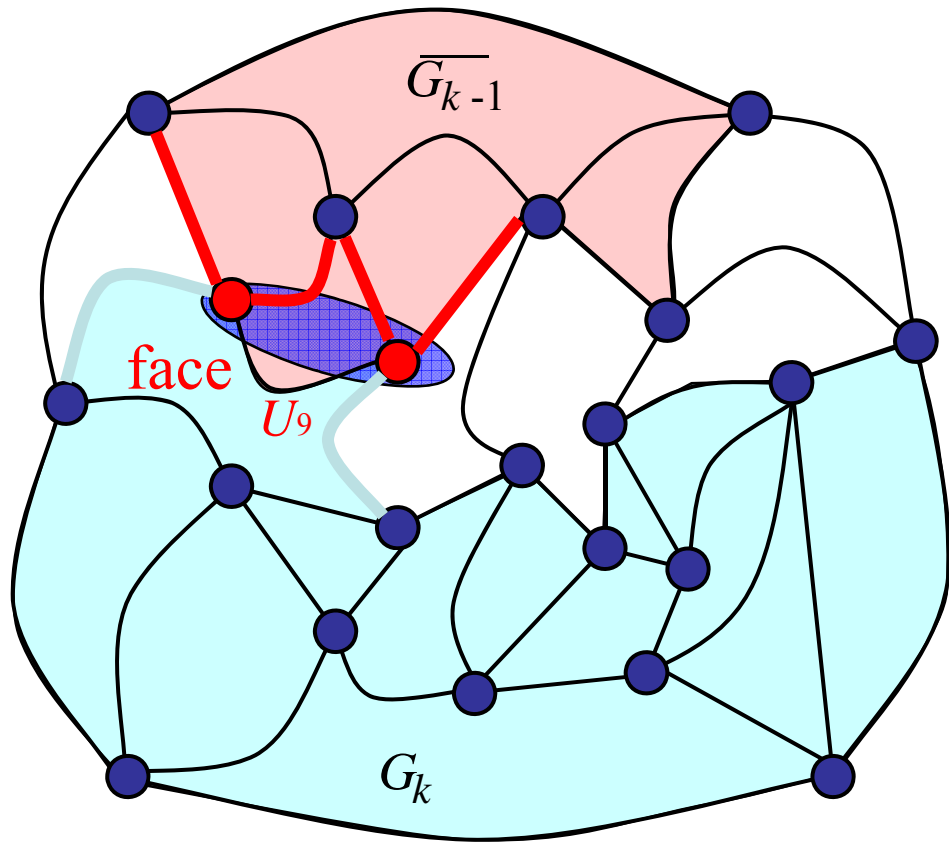
# 4-canonical decomposition [NRN97]

(a generalization of *st*-numbering)



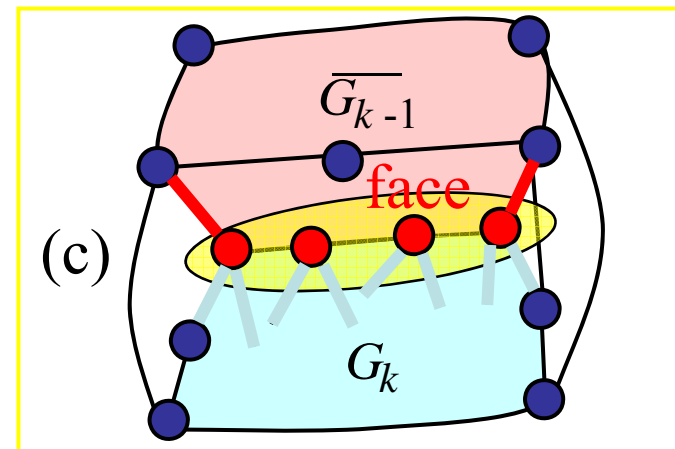
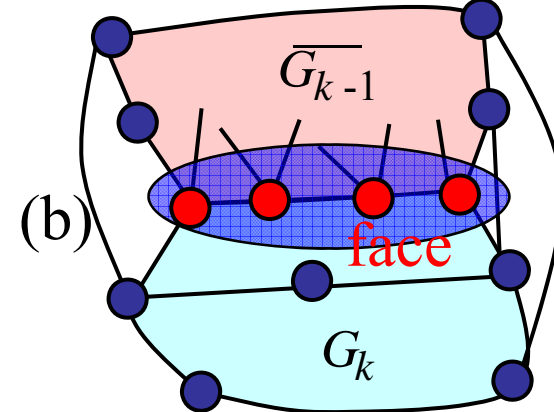
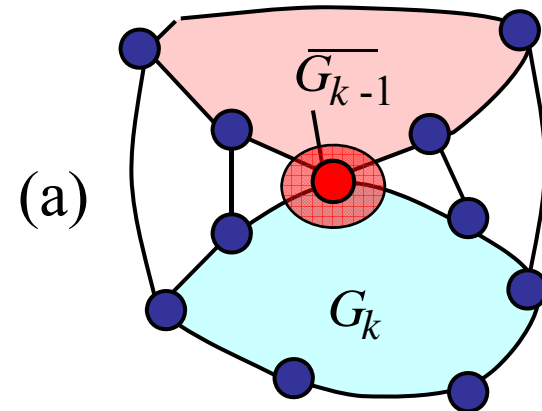
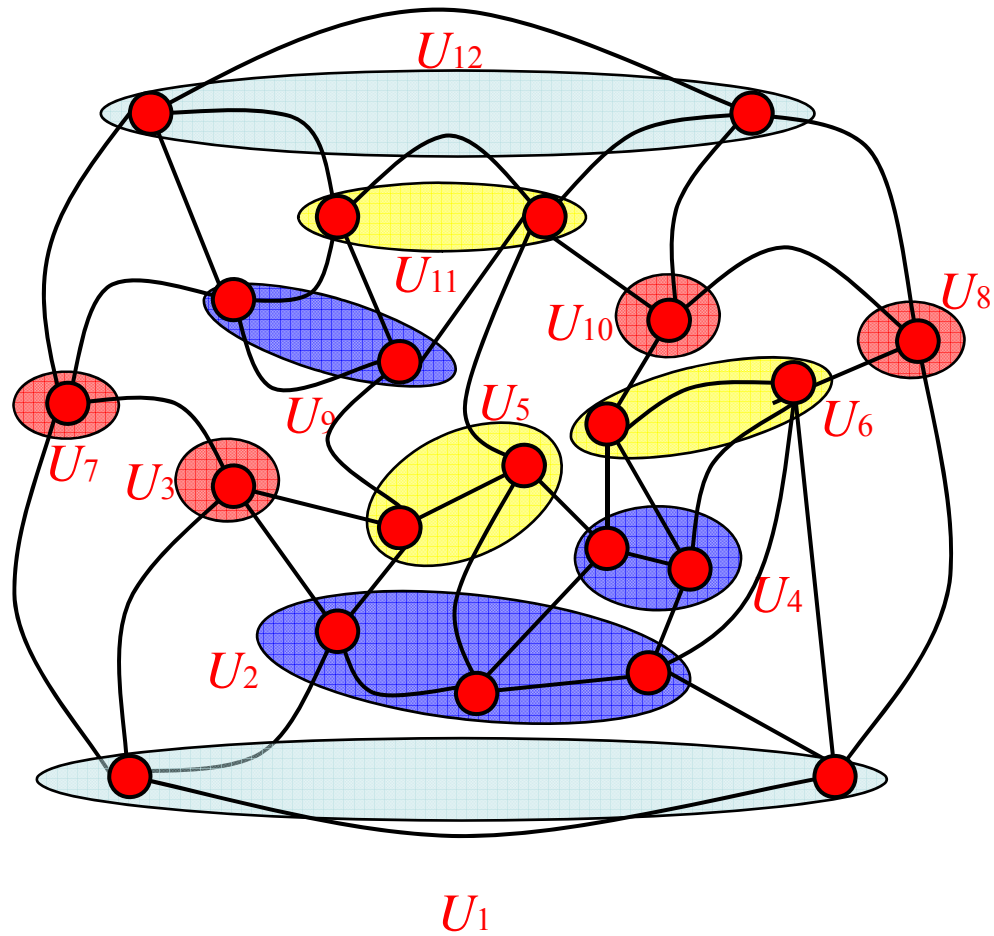
# 4-canonical decomposition [NRN97]

(a generalization of *st*-numbering)



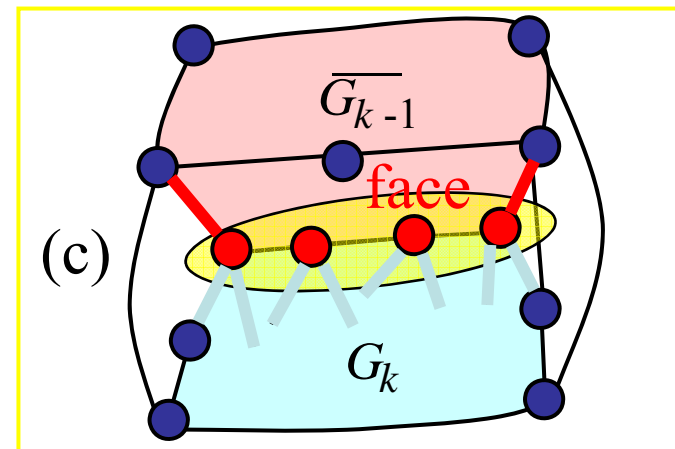
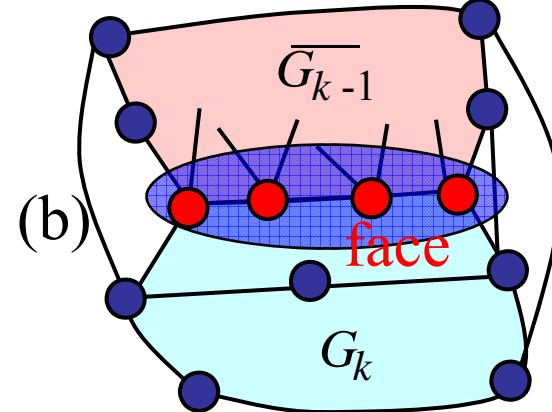
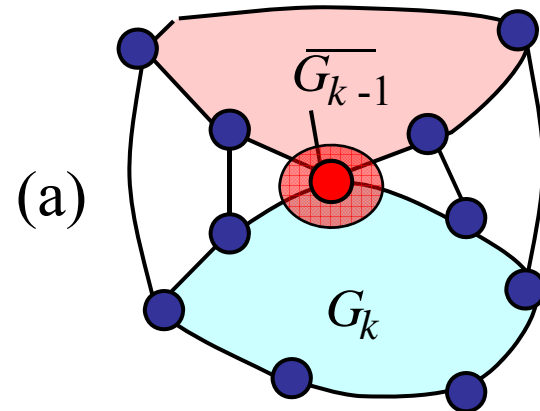
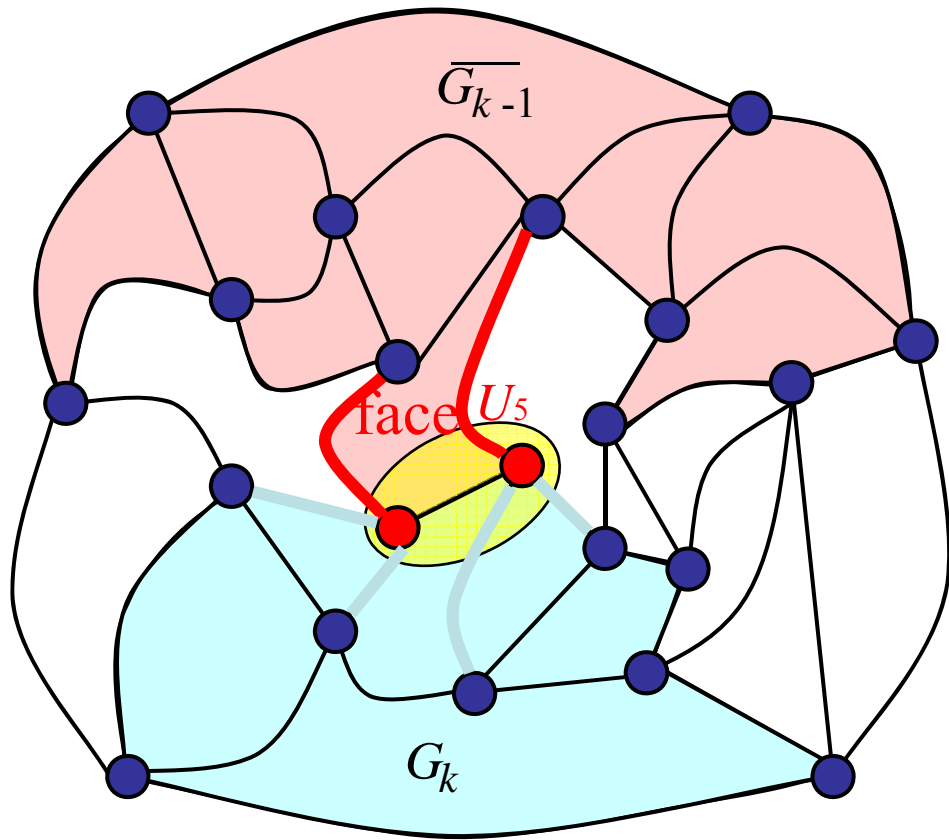
# 4-canonical decomposition [NRN97]

(a generalization of *st*-numbering)



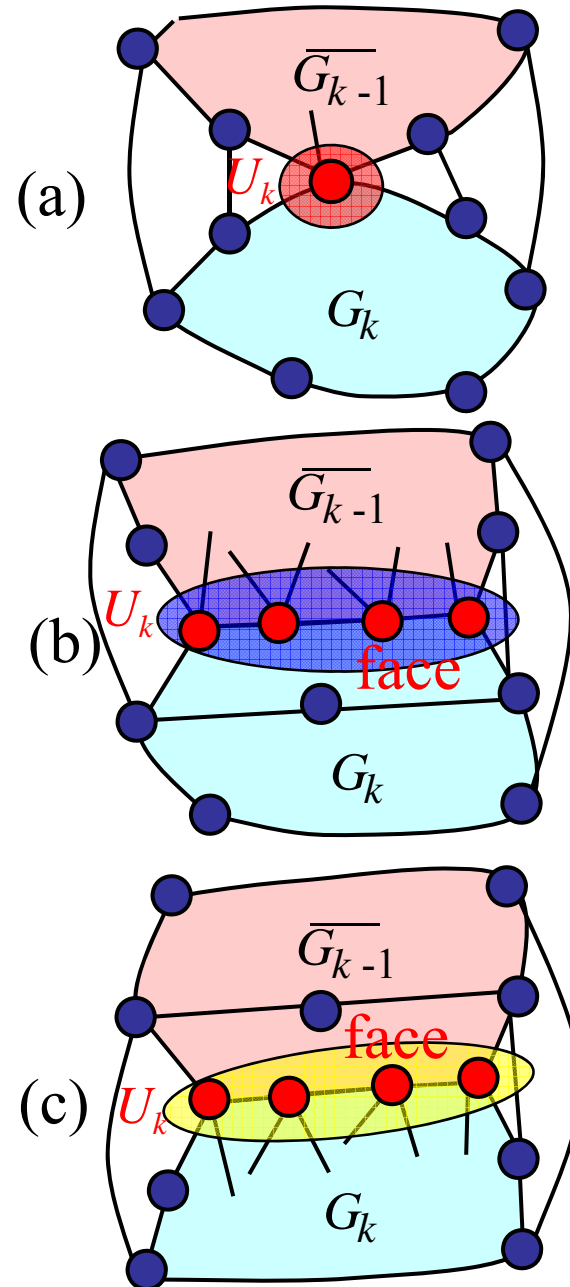
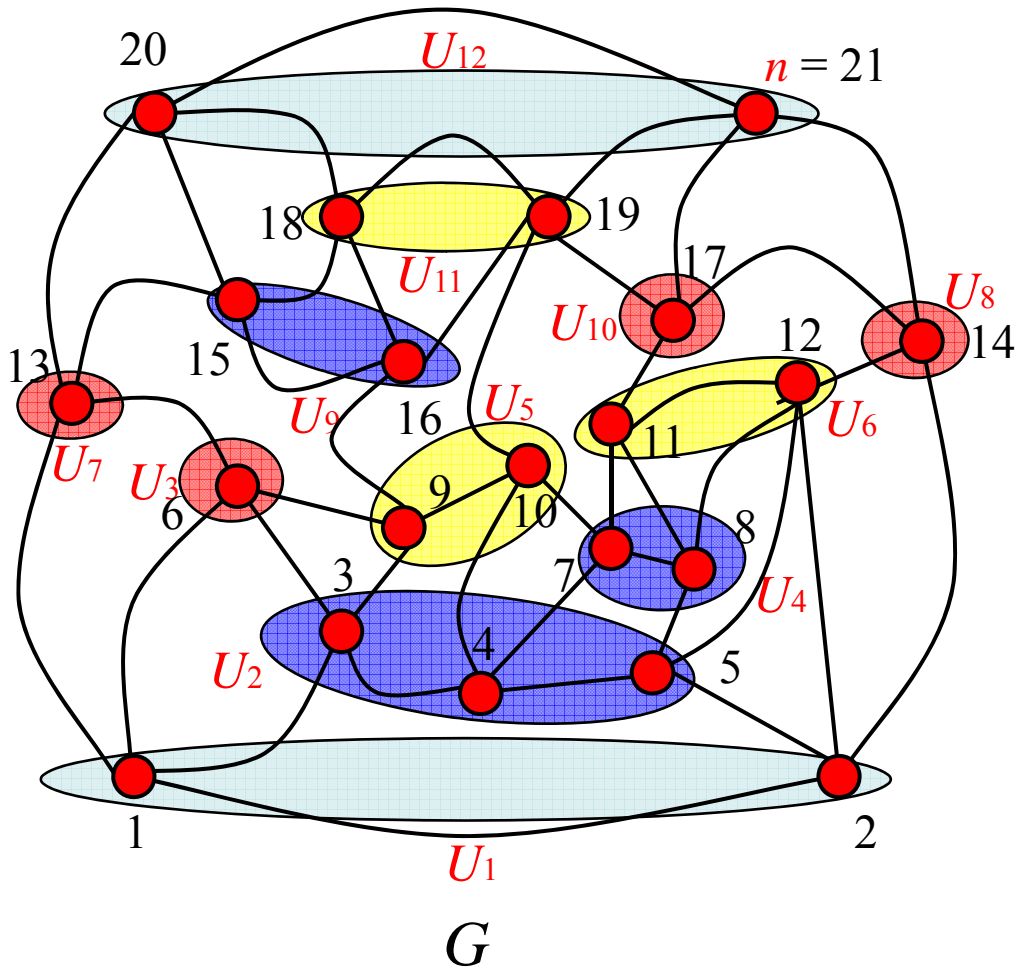
# 4-canonical decomposition [NRN97]

(a generalization of *st*-numbering)



# 4-canonical decomposition [NRN97] $O(n)$

(a generalization of *st*-numbering)



## Conclusions

- *st*-numbering
- Canonical ordering
- 4-canonical ordering
- Canonical decomposition
- 4-canonical decomposition

# Conclusions

## Recent Development

Chiang *et al.*, 2001      Orderly spanning trees

Miura *et al.* 2004

Canonical decomposition, realizer and orderly  
Spanning tree are **equivalent notions**.



Thank You